

INTRODUCTION TO HIGHER ALGEBRA

§1: SYMMETRIC MONOIDAL $\rightarrow$ OPERADS $\rightarrow$ ∞ -LAND

The basic algebraic structure we want to generalise is

commutative monoid: Set M + multiplication $M \times M \rightarrow M$, unit $1 \in M$ s.t.

$$1x = x, xy = yx, x(yz) = (xy)z \quad \forall x, y, z \in M$$

$\leadsto$ For categories this is a symmetric monoidal category: Category C + $1 \in C$ unit

When working with categories its unnatural to ask for $+ \otimes: C \times C \rightarrow C$

$X \otimes (Y \otimes Z) = (X \otimes Y) \otimes Z$, instead we want this structure to be given by extra data in form isomorphisms

$$\alpha_X: 1 \otimes X \simeq X$$

$$\beta_{X,Y}: X \otimes Y \simeq Y \otimes X$$

Just called monoidal if we do not assume this

$$\gamma_{X,Y,Z}: X \otimes (Y \otimes Z) \simeq (X \otimes Y) \otimes Z$$

+ coherence data!

It is clear that this approach can not be used to generalize to ∞ -categories. Lets spell out how to get an equivalent description of symmetric monoidal: The idea is that instead of giving the bifunctor $\otimes: C \times C \rightarrow C$, we instead for each n -tuple $C_1, \dots, C_n$ and $d \in C$ we specify the collection of maps

$$C_1 \otimes \dots \otimes C_n \rightarrow d \quad + \text{ composition data (and coherence)}$$

This is done by colored operads (what I called multicategories in an earlier talk):

Def: A coloured operad $\mathcal{O}$ consists of

- A set of objects $\text{ob } \mathcal{O}$ (sometimes called the colours)
- $\forall x_0, \dots, x_n, y$ in $\mathcal{O}$: A set of multimorphisms $\text{Mul}_{\mathcal{O}}(x_0, \dots, x_n; y)$
- Composition: Given

$$(x_1, \dots, x_{i_1}) \rightarrow y_1, \dots, (x_{i_{n-1}}, \dots, x_{i_n}) \rightarrow y_n$$

can compose this with $(y_1, \dots, y_n) \rightarrow z$ to obtain

$$(x_1, \dots, x_{i_n}) \rightarrow z$$

- Unit: $\text{id}_x: \{x\} \rightarrow x$

s.t. the composition law is unital and associative

Rem: Every colored operad $\mathcal{O}$ has an underlying category by setting:

- objects = $\text{ob } \mathcal{O}$
- $\text{Hom}(x, y) = \text{Mul}_{\mathcal{O}}(\{x\}, y)$

~ So can view a coloured operad as a category + extra data in form of the collection of multimorphisms.

Rem: Given a (symmetric) monoidal category $(C, \otimes)$ we can obtain a coloured operad $C^\otimes$ w. underlying category C by assigning

$$\text{Mul}_C(X_0, \dots, X_n; Y) = \text{Hom}_C(X_0 \otimes \dots \otimes X_n, Y).$$

We can recover the symmetric monoidal structure of $C^\otimes$ (up to canonical isomorphism) by Yoneda's Lemma. E.g. the tensor product $X \otimes Y$ is characterized by the fact that it corepresents the functor $Z \mapsto \text{Mul}_C(X, Y; Z)$.

~> Can consider symmetric monoidal categories as a special case of coloured operads.

To identify which assumptions it is on a coloured operad that makes it into a symmetric monoidal.

First recall:

Def: $\text{Fin}_* = \{\langle n \rangle\}$ pointed sets

$f: \langle n \rangle \rightarrow \langle m \rangle$ is inert if it takes some point to the basepoint and injective (isomorphic) on the rest.

$p: \langle n \rangle \rightarrow \langle 1 \rangle$ takes $j \mapsto 0$ if $j \neq i$, $i \mapsto i$, p is active if $p^{-1}(0) = \emptyset$

Construction: Let $\mathcal{O}$ be a coloured operad ~> $\mathcal{O}^\otimes$ category:

2.1.1.7.

- Objects = sequences of objects of $\mathcal{O}$ $X_0, \dots, X_n$
- $\{X_i\}_{1 \leq i \leq m} \rightarrow \{Y_j\}_{1 \leq j \leq n}$ is given by a map $\alpha: \langle m \rangle \rightarrow \langle n \rangle$ in Fin_* together with a collection of multimorphisms

$$\{\varphi_j \in \text{Mul}_{\mathcal{O}}(\{X_i\}_{i \in \alpha^{-1}\{j\}}, Y_j)\}_{0 \leq j \leq n}$$

in $\mathcal{O}$.

- Composition of morphisms in $\mathcal{O}^\otimes$ is determined by composition laws on Fin_* and on $\mathcal{O}$.

By construction, $\mathcal{O}^\otimes$ comes equipped with a forgetful functor

$$\begin{aligned} \pi: \mathcal{O}^\otimes &\rightarrow \text{Fin}_* \\ (X_0, \dots, X_n) &\mapsto \langle n \rangle. \end{aligned}$$

Using π we can reconstruct the operad structure:

- Write $\mathcal{O}_{\langle n \rangle}^\otimes = \pi^{-1}\{\langle n \rangle\}$

- $p_i^1: \langle n \rangle \rightarrow \langle 1 \rangle \rightsquigarrow p_i^1: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}_{\langle 1 \rangle}^\bullet \cong \mathcal{O}$ ← underlying category of $\mathcal{O}$
which induces equivalences: $\mathcal{O}_{\langle n \rangle}^\bullet \cong \mathcal{O}^{x_n}$.
 $(x_0, -, x_n) \mapsto \overline{x}$.

- $\text{Mul}_{\mathcal{O}}(x_0, -, x_n; Y) \Leftrightarrow \left\{ f: \overline{x} \rightarrow Y \text{ in } \mathcal{O}^\bullet \text{ s.t. } \pi(f): \langle n \rangle \rightarrow \langle 1 \rangle \text{ satisfies } \pi(f)^{-1}(0) = \{0\} \right\}$
↳ $\pi: \mathcal{O}^\bullet \rightarrow \text{Fin}_*$ determines $\text{Mul}_{\mathcal{O}}(x_0, -, x_n; Y)$ in the colored operad $\mathcal{O}$
- Can show that composition law for morphisms in $\mathcal{O}$ can be recovered from the one in $\mathcal{O}^\bullet$.

↳ Can think of a coloured operad as an ordinary category $\mathcal{O}^\bullet$ together with a forgetful functor $\pi: \mathcal{O}^\bullet \rightarrow \text{Fin}_*$ s.t. $\mathcal{O}_{\langle n \rangle}^\bullet \cong (\mathcal{O}_{\langle 1 \rangle}^\bullet)^{x_n}$.
It turns out that if π is further an "op-fibration" then this is exactly the symmetric monoidal categories.

So this is what we generalise to ∞ -land.

Def: An ∞ -operad is a functor of ∞ -categories $p: \mathcal{O}^\bullet \rightarrow N(\text{Fin}_*)$ s.t.

- 1) coCart. lift of every inert $f: \langle n \rangle \rightarrow \langle m \rangle$
↳ In particular it induces $f_!: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}_{\langle m \rangle}^\bullet$.
 - 2) For each $n \geq 0$, the functors $\{p_i^2: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}\}_{1 \leq i \leq n}$ determines an equivalence of categories $\mathcal{O}^{x_n} \cong \mathcal{O}_{\langle n \rangle}^\bullet$.
- + ...

So we in particular get that $(x_0, -, x_n) \in \mathcal{O}^{x_n}$ corresponds to an object of $\mathcal{O}_{\langle n \rangle}^\bullet$.

Def: A symmetric monoidal ∞ -category is an ∞ -operad $\mathcal{C}^\bullet \rightarrow \text{Fin}_*$ which is also coCartesian fibration.

Let's understand why this gives the desired structure:

Note that: $\langle n \rangle \rightarrow \langle m \rangle$ in $\text{Fin}_* \rightsquigarrow \mathcal{C}_{\langle n \rangle}^\bullet \rightarrow \mathcal{C}_{\langle m \rangle}^\bullet$

so the active morphisms $\langle 0 \rangle \rightarrow \langle 1 \rangle, \langle 2 \rangle \rightarrow \langle 1 \rangle$

$\rightsquigarrow$ $\underbrace{\Delta^\bullet \rightarrow \mathcal{C}}_{\text{An object we denote by } 1}, \underbrace{\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}}_{\text{Our tensor}}$

The reason we went through all this work introducing ∞ -operads is because

they are exactly what we need to get nice algebraic structures, and symmetric monoidal categories is just one of many 'monoidal'-ish categories.

In general:

Def: We say that $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$ exhibits $\mathcal{C}$ as $\mathcal{O}$ -monoidal ω -category if

- 1) p is coCartesian
- 2) $\mathcal{O}^\otimes \xrightarrow{q} \text{Fin}_*$ is an ω -operad and $\mathcal{C}^\otimes \xrightarrow{p} \mathcal{O}^\otimes \xrightarrow{q} \text{Fin}_*$ exhibits $\mathcal{C}^\otimes$ as an ω -operad.

Ex: Symmetric monoidal $\leftrightarrow$ Fin_* -monoidal.

Rem: If $\mathcal{C}$ is $\mathcal{O}$ -monoidal by $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$, then any

$$X \in \mathcal{O}^\otimes_{\langle n \rangle} \leftrightarrow \{X_0, \dots, X_n\} \in \mathcal{O}$$

Given

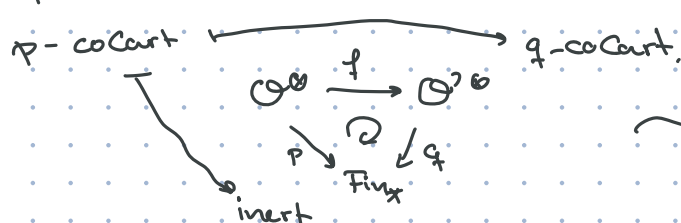
$$\text{f} \in \text{Mul}_{\mathcal{O}}(X_0, \dots, X_n; Y) \rightsquigarrow \prod_{0 \leq i \leq n} \mathcal{C}_{X_i} \xrightarrow{\text{f}} \mathcal{C}_Y$$

$\uparrow$
 fiber over x

It can be shown that one of the key elements is that p in particular is a morphism of ω -operads:

ALGEBRA OBJECTS

Def: A morphism of ω -operads :

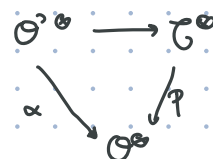


$$\text{Alg}_{\mathcal{O}}(\mathcal{O}') \subseteq \text{Fun}_{\text{Fin}_*}(\mathcal{O}^\otimes, \mathcal{O}'^\otimes)$$

i.e. morphism + fibration on underlines

Def: Assume $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$ be a fibration of ω -operads and given $\alpha: \mathcal{O}'^\otimes \rightarrow \mathcal{O}^\otimes$

$$\text{Alg}_{\mathcal{O}'/\mathcal{O}}(\mathcal{C}) \subseteq \text{Fun}_{\mathcal{O}^\otimes}(\mathcal{O}'^\otimes, \mathcal{C}^\otimes)$$



Equivalently:

$$\text{Alg}_{\mathcal{O}'/\mathcal{O}}(\mathcal{C}) = \text{fib. of } \text{Alg}_{\mathcal{O}}(\mathcal{C}) \xrightarrow{p} \text{Alg}_{\mathcal{O}}(\mathcal{O}) \text{ at } \alpha$$

$$\mathcal{O}'^\otimes \rightarrow \mathcal{C}^\otimes \mapsto \mathcal{O}'^\otimes \rightarrow \mathcal{C}^\otimes \xrightarrow{p} \mathcal{O}^\otimes$$

Case $\mathcal{O}^\otimes = \mathcal{O}'^\otimes, \alpha = \text{id}$: $\text{Alg}_{\mathcal{O}/\mathcal{O}}(\mathcal{C}) =: \text{Alg}_{/\mathcal{O}}(\mathcal{C})$

Case $\Theta^{\otimes} = \Theta^{\otimes} = \text{Fin}_*$:

$$\text{CAlg}(\mathcal{C}) := \text{Alg}_{\Theta/\mathcal{C}}(\mathcal{C}) = \text{Alg}_{/\mathcal{C}}(\mathcal{C})$$

$$\begin{array}{ccc} \text{Fin}_* & \xrightarrow{\quad} & \mathcal{C}^{\otimes} \\ \parallel & \searrow & \\ \text{Fin}_* & & \end{array}$$

It's specific cases which gives us algebras!

Base case is associative algebra

Classically: An associative algebra object in a monoidal category $\mathcal{C}$ is an object $A \in \mathcal{C}$ + unit map

$e: 1 \rightarrow A$ and multiplication $m: A \times A \rightarrow A$ s.t.

$$\begin{array}{ccc} 1 \otimes A & \xrightarrow{\text{e} \otimes \text{id}} & A \otimes A \\ & \searrow & \swarrow m \\ & A & \end{array} \quad \begin{array}{ccc} A \otimes 1 & \xrightarrow{\quad} & A \otimes A \\ & \searrow & \swarrow m \\ & A & \end{array}$$

$$\begin{array}{ccc} A \otimes A \otimes A & \xrightarrow{\text{m} \otimes \text{id}} & A \otimes A \\ \downarrow \text{id} \otimes m & & \downarrow m \\ A \otimes A & \xrightarrow{m} & A \end{array}$$

We will apply the formalism of ∞ -operads to introduce the notion of monoidal ∞ -category, and to each monoidal ∞ -category associate another ∞ -category $\text{Alg}(\mathcal{C})$ of associative algebra objects of $\mathcal{C}$.

Def: Colored operad **Assoc** (associative algebra)

- $\text{ob Assoc} = \{a\}$
- $\text{Mul}_{\text{Assoc}}(\{a\}_{i \in I}, a) = \text{set of linear orderings on } I.$
any finite set.
- + composition...

$\leadsto$ Obtain a category $\text{Assoc}^{\otimes}$ by applying construction 2.1.1.7

Unwinding the def of $\text{Assoc}^{\otimes}$:

- $\text{ob Assoc}^{\otimes} = \text{ob Fin}_*$
- $\langle m \rangle \rightarrow \langle n \rangle$ in $\text{Fin}_* \leadsto \langle m \rangle \rightarrow \langle n \rangle$ in $\text{Assoc}^{\otimes}$ consists of a pair $(\kappa, \{i_1\}_{1 \leq i_1 \leq m})$ where $\kappa: \langle m \rangle \rightarrow \langle n \rangle$ is a map in Fin_* and $\leq_{i_1}$ is a linear ordering on the inverse image $\kappa^{-1}\{i_1\} \subseteq \langle m \rangle$ for $1 \leq i_1 \leq n$

$\leadsto \text{Assoc}^{\otimes} = N(\text{Assoc}^{\otimes}) \in \text{Cat}_{\infty}$

Notation: We write $\text{Assoc} := \text{Assoc}^{\otimes} \times_{\text{Fin}_*} \{\langle 1 \rangle\}$

Note that as a simplicial set, Assoc is isomorphic to the 0-simplex Δ^0 ; But we use the notation Assoc to emphasize the role of the simplicial set as the underlying ∞ -category for the ∞ -operad $\text{Assoc}^\otimes$.

Def: A monoidal ∞ -category is a coCartesian fibration of ∞ -operads $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$.

Def: $\mathcal{C}^\otimes \in \text{Op}^\otimes$ equipped w. a fibration $q: \mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$. Then the ∞ -category of associative algebra objects of $\mathcal{C}$ is $\text{Alg}(\mathcal{C}) := \text{Alg}_{/\text{Assoc}}(\mathcal{C})$ (∞ -operad sections of q)

$$\begin{array}{ccc} \text{Assoc}^\otimes & \xrightarrow{A} & \mathcal{C}^\otimes \\ \text{Id} \swarrow & & \searrow q \\ & \text{Assoc}^\otimes & \end{array}$$

Let's understand what this means: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a monoidal ∞ -category. Then

rem. 4.1.12-15 $\mathcal{C}_{\langle n \rangle}^\otimes \cong \mathcal{C}^\otimes \times_{\text{Assoc}^\otimes} \{\langle n \rangle\} \cong \mathcal{C}^n$, and for every linear ordering on $\{1, \dots, n\}$, the corresponding map $\langle n \rangle \rightarrow \langle 1 \rangle$ in $\text{Assoc}^\otimes$ induces a functor $\mathcal{C}^n \rightarrow \mathcal{C}$. In particular

- $n=0 \leadsto$ unit object $1 \in \mathcal{C}$
- $n=2$, standard ordering on $\{1, 2\} \leadsto \otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$
- Evaluating on $a \in \text{Assoc}^\otimes$ determines a forgetful functor

$$\Theta: \text{Alg}(\mathcal{C}) \rightarrow \mathcal{C}$$

By abuse of notation we often identify $A \in \text{Alg}(\mathcal{C})$ with its image $\Theta(A)$ in $\mathcal{C}$. For each $n \geq 0$ a choice of ordering on $\{1, \dots, n\}$ determines an active morphism $\{a_i\}_{1 \leq i \leq n} \rightarrow a$ in $\text{Assoc}^\otimes$, which induces a morphism $\Theta(A)^{\otimes n} \rightarrow \Theta(A)$ in $\mathcal{C}$. In particular

- $n=2$ and standard ordering of $\{1, 2\}$: $m: \Theta(A) \otimes \Theta(A) \rightarrow \Theta(A)$
- $\hookrightarrow$ Associative and unital up to homotopy

$\leadsto$ In particular, it endows $\Theta(A)$ with the structure of an associative algebra object in $\text{h}\mathcal{C}$.

LEFT & RIGHT MODULES

Classically: $\mathcal{C}$ monoidal category w. unit A , A associative algebra object of $\mathcal{C}$. A left A -module in $\mathcal{C}$ is an object $M \in \mathcal{C}$ + action map $a: A \otimes M \rightarrow M$ s.t.

$$\begin{array}{ccc} A \otimes A \otimes M & \xrightarrow{m \circ id} & A \otimes M \\ \downarrow id \otimes a & & \downarrow a \\ A \otimes M & \xrightarrow{a} & M \end{array}$$

$$\begin{array}{ccc} 1 \otimes M & \xrightarrow{u \circ id} & A \otimes M \\ & \searrow & \swarrow a \\ & M & \end{array}$$

where $m = \text{multiplication}$
 $u = \text{Unit}$ } in A

$\leadsto$ All left A -modules = $\text{LMod}_A(\mathcal{C})$

We wish to introduce a larger $\mathcal{A}\mathcal{O}$ -operad $\text{LM}^\otimes$ which contains $\text{Assoc}^\otimes$. If $A \in \text{Alg}(\mathcal{C})$ we want a left A -module to be a map of $\mathcal{A}\mathcal{O}$ -operads $M: \text{LM}^\otimes \rightarrow \mathcal{C}^\otimes$ s.t. $M|_{\text{Assoc}^\otimes} = A$.

Def: Define a colored operad $\underline{\text{LM}}$ as:

- $\text{ob LM} = \{a, m\}$
- Let $\{X_i\}_{i \in I}$ be a finite collection of objects of LM . Then
 - $\hookrightarrow \text{Mul}_{\text{LM}}(\{X_i\}, a) = \begin{cases} \text{all linear orderings of } I \text{ if all } X_i = a \\ \text{Empty} & \text{otherwise} \end{cases}$
 - $\hookrightarrow \text{Mul}_{\text{LM}}(\{X_i\}, m) = \{ \text{All linear orderings } \{i_1, \dots, i_n\} \text{ of } I \text{ s.t. } X_{i_1} = m \text{ \& } X_{i_j} = a \text{ for } j > 1 \}$

Rem:
4.2.1.2

$a \in \text{LM} \leadsto$ sub-colored operad of LM isomorphic to Assoc

$a, a, \dots, a, m$

Rem:
4.2.1.3

We first in this (1-categorical) case see how this can be used to understand modules.

Assume $\mathcal{C}$ symmetric monoidal, $F: \text{LM} \rightarrow \mathcal{C}$ map of colored operads.

$\leadsto F|_{\text{Assoc}}: \text{Assoc} \rightarrow \mathcal{C} \leadsto$ Associative algebra object $F(a) = A \in \mathcal{C}$. Let $M = F(m) \in \mathcal{C}$.

Then the unique operation $\phi \in \text{Mul}_{\text{LM}}(\{a, m\}, m)$ determines $F(\phi): A \otimes M \rightarrow M$, which exhibits M as a left A -module.

$\leadsto$ So $\begin{array}{l} \text{LM} \rightarrow \mathcal{C} \\ a \mapsto \text{associative algebra} \\ m \mapsto \text{left module over that algebra} \end{array}$

Notation: Apply construction 2.1.1.7 to LM to obtain the category $\text{LM}^\otimes$ from LM . Unwinding this construction we see that

$$1) \text{ ob LM}^\otimes = \{(\langle n \rangle, S) \mid S \subset \langle n \rangle^\otimes\}$$

$$2) (\langle n \rangle, S) \rightarrow (\langle n' \rangle, S') = \alpha: \langle n \rangle \rightarrow \langle n' \rangle \text{ in } \text{Assoc}^\otimes \text{ s.t.}$$

$$i) S \cup \{*\} \xrightarrow{\alpha} S' \cup \{*\}$$

ii) $S' \in S' \Rightarrow \alpha^{-1}\{s'\}$ contains exactly one element of S , and that object is maximal w.r.t. the linear ordering of $\alpha^{-1}\{s'\}$.

Rem:

$$\begin{array}{ccc} LM & & LM^\otimes \\ a & \longleftarrow & (\langle 1 \rangle, \emptyset) \\ m & \longleftrightarrow & (\langle 1 \rangle, \langle 1 \rangle^\circ) \end{array}$$

Def: $LM^\otimes = N(LM^\otimes)$. This is an ∞ -operad via the forgetful map $LM^\otimes \rightarrow \text{Fin}_*$

Rem: The underlying ∞ -category LM of $LM^\otimes$ is isomorphic to the discrete simplicial set $\Delta[1]$ w. two vertices, corresponding to $a, m \in LM$.

Rem: $\text{Assoc} \hookrightarrow LM \rightsquigarrow \text{Assoc}^\otimes \hookrightarrow LM^\otimes$, which is an isomorphism from $\text{Assoc}^\otimes$ onto the full subcategory of $LM^\otimes$ spanned by objects of the form $(\langle n \rangle, \emptyset)$.

Notation: Let $\mathcal{C}^\otimes \rightarrow LM^\otimes$ be a fibration of ∞ -operads. Write

$$\mathcal{C}_a^\otimes := \mathcal{C}^\otimes \times_{LM^\otimes} \text{Assoc}^\otimes, \quad \text{Underlying } \infty\text{-category } \mathcal{C}_a = \mathcal{C}^\otimes \times_{LM^\otimes} \{a\}$$

$$\mathcal{C}_m^\otimes = \mathcal{C}^\otimes \times_{LM^\otimes} \{m\}$$

Def: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a fibration of ∞ -operads, $q: \mathcal{C}^\otimes \rightarrow LM$ fibration of ∞ -operads s.t. $\mathcal{C}_a^\otimes \simeq \mathcal{C}^\otimes$. Write $\mathcal{C}_m^\otimes = M_0$ ($\in \text{Cat}_\infty$) (Normally we say q exhibits M_0 as weakly enriched over $\mathcal{C}^\otimes$)

- $L\text{Mod}(M) := \text{Alg}_{/LM}(\mathcal{C})$ ∞ -category of left module objects of M_0 $\begin{array}{c} LM \rightarrow \mathcal{C} \\ \downarrow \\ LM \end{array}$
- Composition w. $\text{Assoc}^\otimes \hookrightarrow LM^\otimes$ determines a categorical fibration

$$L\text{Mod}(M_0) = \text{Alg}_{/LM}(\mathcal{C}) \xrightarrow{\circ \text{incl}} \text{Alg}_{/\text{Assoc}}(\mathcal{C}) = \text{Alg}(\mathcal{C})$$

$$\rightsquigarrow L\text{Mod}_A(M) = L\text{Mod}(M) \times_{\text{Alg}(\mathcal{C})} \{A\} \quad \infty\text{-category of left } A\text{-modules objects of } M_0.$$

$\rightsquigarrow$ We think of $L\text{Mod}(\mathcal{C}_m)$ as given by pairs (A, M) where A is an associative algebra object of $\mathcal{C}_a$ and M is a left A -module in $\mathcal{C}_m$.

Ex: $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ fibration of ∞ -operads, $\mathcal{C}^\otimes = \mathcal{C}^\otimes \times_{\text{Assoc}^\otimes} LM^\otimes$. Then " $\mathcal{C}^\otimes$ exhibits $\mathcal{C}$ as weakly enriched over $\mathcal{C}^\otimes$ " $\rightsquigarrow$ Can consider $L\text{Mod}(\mathcal{C}) = \text{Alg}_{LM/\text{Assoc}}(\mathcal{C})$

$$\begin{array}{ccc} LM^\otimes & \xrightarrow{\circ} & \mathcal{C}^\otimes \\ & \searrow & \downarrow \\ & & \text{Assoc}^\otimes \end{array}$$

Ex: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a monoidal ∞ -category. Then

$$L\text{Mod}(\mathcal{C}) = \text{Alg}_{LM/\text{Assoc}}(\mathcal{C}) = LM^\otimes \xrightarrow{F} \mathcal{C}^\otimes$$

$\downarrow$
 $\text{Assoc}^\otimes$

$\rightsquigarrow F|_{\text{Assoc}} \in \text{Alg}(\mathcal{C})$, which we identify with its underlying object $F(a) = A \in \mathcal{C}$.

$\rightsquigarrow$ Also have $F(m) = M \in \mathcal{C}$

The unique operation $q \in \text{Mod}_{LM}(\{a, m\}, m)$ determines a map

$$a: A \otimes M \rightarrow M \quad \text{in } \mathcal{C}$$

which is well-defined up to homotopy.

Since F is defined on all of $LM^\otimes$, we get that this action map is compatible with the associative multiplication on A , up to coherent homotopy. In particular, if

$$\begin{aligned} m: A \otimes A &\rightarrow A && \text{multiplication on } A \\ u: 1 &\rightarrow A && \text{unit map} \end{aligned}$$

$$\begin{array}{ccc} A \otimes A \otimes M & \xrightarrow{m \otimes \text{id}} & A \otimes M \\ \downarrow \text{id} \otimes a & & \downarrow a \\ A \otimes M & \xrightarrow{a} & M \end{array} \quad \begin{array}{ccc} 1 \otimes M & \xrightarrow{u \otimes \text{id}} & A \otimes M \\ & \searrow & \swarrow a \\ & M & \end{array}$$

commutes up to homotopy. of ∞ -ops.

Rem: If $\mathcal{C}^\otimes \xrightarrow{q} LM^\otimes$ is a cocartesian fibration, then the induced map

$\mathcal{C}_a^\otimes \rightarrow \text{Assoc}^\otimes$ is also a cocartesian fibration of ∞ -operads.

We further see that by straightening, q is classified by a map $\chi: LM^\otimes \rightarrow \text{Cat}_\infty$, which is an " LM -monoid object of Cat_∞ ". We have

$$\chi \in \text{Mon}_{LM}(\text{Cat}_\infty) \simeq \text{Alg}_{LM}(\text{Cat}_\infty) = \text{LMod}(\text{Cat}_\infty)$$

More informally: q can be thought of as giving an associative algebra $\mathcal{C}_a$ in Cat_∞ together with a left module $\mathcal{C}_m$ over $\mathcal{C}_a$. In particular q determines an action map

$$\otimes: \mathcal{C}_a \times \mathcal{C}_m \rightarrow \mathcal{C}_m$$

which is well-defined up to homotopy - and compatible with the monoidal structure on $\mathcal{C}_a$.