

Paul 05.10.22: Introduction

Motivation

X : a topological space $\leadsto \pi_0 X, x \in X, \pi_1(X, x)$

$$\pi_{\leq 1}(X, x) = \text{cat} \begin{cases} \text{ob} = \text{pts in } X \\ \text{morp} = \text{homotopy classes of paths in } X \end{cases}$$

$$\text{Hom}_{\pi_1}(x, x) = \pi_1(X, x)$$

Q: What sort of "gadget" contains the data of a homotopy theory? — ∞ -categories

- All (small) categories w. functors and nat. trans form a 2Cat.
- Convenient way to do homotopy theory!

Def: A poset \mathcal{P} is a set with a partial order

Every poset \mathcal{P} determines a category:

$$\begin{aligned} \text{ob}(\mathcal{P}) &= \mathcal{P} \\ \text{Hom}_{\mathcal{P}}(x, y) &= \begin{cases} \emptyset & x \not\leq y \\ x & x \leq y \end{cases} \end{aligned}$$

Ex: $[n] = \{0 \leq 1 \leq \dots \leq n\}, n \geq 0$

$$\Delta = \begin{cases} \text{ob} = \{[n]\} \\ \text{Hom}_{\Delta}([n], [m]) = \text{Fun}([n], [m]) \end{cases} \quad \text{order preserving}$$

Distinguished maps:

$$d_i: [n] \rightarrow [n+1] \quad \text{injective and misses } i$$

$$s_j: [n] \rightarrow [n+1] \quad \text{surjective and } |s_j^{-1}(i)|=2$$

Def: A simplicial set is a functor $\Delta^{\text{op}} \rightarrow \text{Set}$

$X_0 := X([0])$ the 'objects' of X

$X_1 := X([1])$ the 'morphisms' of X

Ex: Δ^n the standard n -simplex $:= \text{Hom}_{\Delta}(\cdot, [n])$

Prop: $X_n \cong \text{Hom}_{\text{Set}}(\Delta^n, X)$

Def: Topological n -simplex $\Delta^n_{\text{top}} = \{(x_0, \dots, x_n) \in \mathbb{R}^{n+1} \mid x_i \geq 0, \sum x_i = 1\}$

Def: Singular simplicial set of a space X is

$$\text{Sing}_n(X) = \text{Hom}_{\text{Top}}(\Delta^n_{\text{top}}, X)$$

Def: Geometric realization of a sSet X :

$$|X| = \bigsqcup X_n \times \Delta^n_{\text{top}} / \sim \text{gluing along faces \& degeneracy maps}$$



This is a topological space (CW-complex).

Prop: There exists an adjunction

$$|\cdot| : \text{sSet} \rightleftarrows \text{Top} : \text{Sing}_\bullet$$

\leadsto Have units & counits

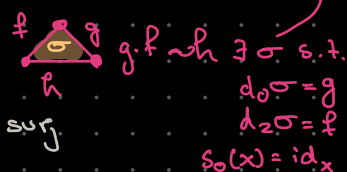
$\hookrightarrow |\text{Sing}_\bullet(X)| \rightarrow X$ weak htpy eq. - CW-approximation

Def: \mathcal{C} a Cat, the nerve $N_\bullet \mathcal{C} : [n] = \text{Fun}([n], \mathcal{C})$. $N : \text{Cat} \rightarrow \text{sSet}$ fully faithful

Def: The homotopy category hX of a simplicial set X :
 $\text{ob } hX = X_0$

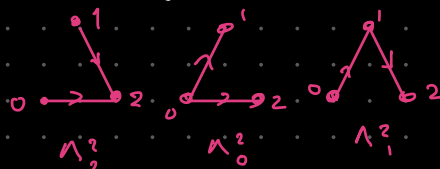
$\text{Hom}_{hX} =$ free compositions of composable morphisms in X_1 / \sim

Punch line: $h \dashv N$, $\mathcal{C} \cong hN\mathcal{C}$



Construction: $(\Delta^n)_2 = \text{maps } [2] \rightarrow [n] \text{ that is not surj.}$

$\Delta^n_j :=$ The j^{th} horn = maps $[k] \rightarrow [n]$ s.t. $\exists i \in [n] \setminus j$ not in the image



$(I^n)_2 := \text{Spine} \subseteq (\Delta^n)_2$ maps $[k] \rightarrow [n]$ whose image is

$$\{j, j+1, j+2\}$$

Def: A Kan complex is a simplicial set with the horn extension property for all horns:

$$\begin{array}{ccc} \bigwedge_{i=1}^n & \xrightarrow{\forall f} & X \\ \downarrow & \nearrow \exists & \\ \Delta^n & & \end{array} \quad \forall n \geq 2, \forall 0 \leq i \leq n$$

Ex: $X \in \text{Top} \leadsto \text{Sing}(X) \in \text{Kan}$ follows by the adjunction $1 \dashv \text{Sing}$ and $1 \dashv \text{Sing} \Rightarrow 1 \dashv \text{Sing}$ def. retracts.

Thm: $N.G$ is a Kan complex $\Leftrightarrow G$ is a groupoid (all morphisms are invertible)

Triv: G a category $\leadsto BG = |N.G|$ is the classifying space of G

\hookrightarrow Book recommendation: "Classifying spaces & classifying topos" - Mac Lane

G group, model it as BG - obj $*$, $\text{Hom}_{BG}(*, *) \cong G$

$|BG|$ = Classifying space of G (in the classical sense)

\hookrightarrow What does this classify for general G ?

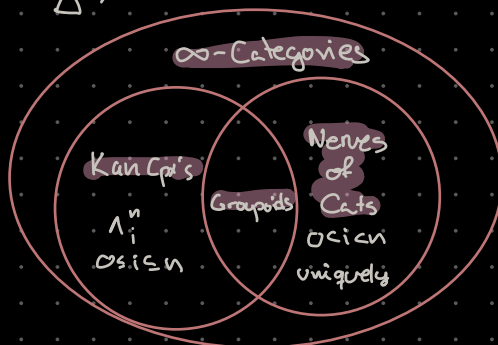
Def: A simplicial set is a cosmos if it has the spine extension property, i.e.

$$\begin{array}{ccc} I^n & \xrightarrow{\forall f} & X \\ \downarrow & \nearrow \exists & \\ \Delta^n & & \end{array} \quad \text{Weaker than } \infty\text{-cat requires}$$

Prop: In $N.G$ there are unique inner horn extensions

Def: An ∞ -category is a simplicial set G with the horn extension property for all inner horns, i.e.

$$\begin{array}{ccc} \bigwedge_{i=1}^n & \xrightarrow{\forall f} & X \\ \downarrow & \nearrow \exists & \\ \Delta^n & & \end{array} \quad \forall n \geq 2, \forall 0 < i < n$$



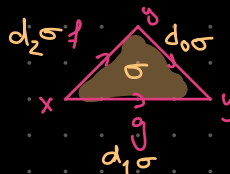
Prop:

- $X \in \text{Top}; \text{Sing. } X \in \text{Cat}_\infty$
- $\mathcal{C} \in \text{Cat}; N.\mathcal{C} \in \text{Cat}_\infty$
- $\mathcal{C} \in \text{Cat}_2; N^{\text{Deskin}} \mathcal{C} \in \text{Cat}_\infty$
- $\mathcal{C} \in \text{Cat}_{\text{Kan}}; N^{\text{hc}} \mathcal{C} \in \text{Cat}_\infty$ ($\mathcal{C} \in \text{Cat}_2$ is only a composer.)

Def:

Say $f, g \in X_1$ are equivalent $f \sim g$ if there exists $\sigma \in X_2$ s.t.

$$\begin{aligned} d_2 \sigma &= f \\ d_1 \sigma &= g \\ d_0 \sigma &= \text{id}_y = S_0(y) \end{aligned}$$



Prop:

If X is a composer & has 2 and/or 3 inner horn extensions, then this is an equivalence then this is an equivalence relation.

Or if X is an ∞ -category!

Def:

$\pi X = 0$ simplices & 1 simplices / \sim equivalent

Prop:

$\pi X \cong \pi X$

Def:

$X \in \text{Cat}_\infty$, call $f \in X_1$ an equivalence if $f \in \pi X$ is an isomorphism.

An ∞ -groupoid is an ∞ -cat where all 1-morphisms are all invertible.

Thm:

Kan complex $X \iff \infty$ -Groupoid

2. Daniel Marlowe - Small object argument

§1: Small object argument

Notation: $\Delta_n = \{\Delta_n^j \hookrightarrow \Delta^n \mid n \geq 1\}$

Inner: $0 < j < n$

Left: $0 < j < n$

Right: $0 < j < n$

Given $\mathcal{S} \subseteq \text{ArSet}$ denote by

$$\mathcal{X}_L(\mathcal{S}) = \{g \text{ w. LLP wrt. } s \in \mathcal{S}\}$$

$$\mathcal{X}_R(\mathcal{S}) = \{g \text{ w. RLP wrt. } s \in \mathcal{S}\}$$

$$\mathcal{X}(\mathcal{S}) := \mathcal{X}_L(\mathcal{X}_R(\mathcal{S}))$$

Def. $\mathcal{S} \subseteq \text{Mor}(\text{sSet})$ is saturated if it is closed under

- i) pushouts
- (ii) arbitrary coproducts) If working w. not just \mathbb{N} -indexed colims, this will follow from i.v.
- iii) Retract
- iv) \mathbb{N} -indexed colimits

$$\begin{array}{c} \text{in } \mathcal{S} \\ x_0 \rightarrow x_1 \rightarrow \dots \rightarrow \text{colim } x_i \\ \Rightarrow \underbrace{\hspace{10em}}_{\text{this is in } \mathcal{S}} \end{array}$$

Ex: • All maps

• Cofibrations

• Monomorphisms

Def $\bar{\mathcal{S}}$ saturated closure

Lem: $\forall \mathcal{S} : \mathcal{X}_L(\mathcal{S})$ is saturated

Rem: $\bar{\mathcal{S}} = \mathcal{X}(\mathcal{S})$

Prop | Quillen | Given $\mathcal{S} = \{A_i \xrightarrow{f_i} B_i\} \text{ s.t.}$

* A_i have $< \infty$ non-degenerate simplices

then for any map of sSets $f: X \rightarrow Y$, \exists

$$\begin{array}{ccc} x & \xrightarrow{s \in \bar{S}} & z \\ & \searrow f & \downarrow p \in \mathcal{K}_R(S) \\ & & y \end{array}$$

PF: Consider the family $\Theta_S = \left\{ \begin{array}{c} A_i \rightarrow x \\ \downarrow f_i \in S \\ B_i \rightarrow y \end{array} \right\}$. Look at

$$\begin{array}{ccc} \coprod_{\Theta_S} A_i & \longrightarrow & x \\ \downarrow \sum f_i & \searrow & \downarrow s_i \in \bar{S} \\ \coprod_{\Theta_S} B_i & \longrightarrow & E^p(f) \xrightarrow{\exists!} y \end{array}$$

f

Iterate $E^p(f) \rightarrow y, \dots$, produce a sequence

$$\begin{array}{c} x \rightarrow E^p(f) \rightarrow \dots \rightarrow y \\ \searrow s \in \bar{S} \quad \nearrow p \\ \quad E^w(f) \end{array}$$

\leadsto Consider

$$\begin{array}{ccc} A_i & \longrightarrow & E^w(f) \\ \downarrow f_i & & \downarrow p \\ B_i & \longrightarrow & y \end{array} \quad \begin{array}{c} \text{colim} \text{Hom}(A_i, E^p(f)) \\ \downarrow \\ \text{Hom}(A_i, E^w(f)) \end{array}$$

$\Rightarrow \exists N$ s.t. we get a factorization

$$\begin{array}{ccccc} A_i & \longrightarrow & E^N(f) & \xrightarrow{c_N} & E^w(f) \\ \downarrow f_i & & \downarrow & \nearrow c_{N+1} & \downarrow p \\ B_i & \longrightarrow & E^{N+1}(f) & \longrightarrow & y \end{array} \Rightarrow p \in \mathcal{K}_R(S)$$

Cor: For S satisfying (*): $\bar{S} = \mathcal{K}(S)$.

PF: $\bar{S} \subseteq \mathcal{K}(S)$ ✓

To see $\mathcal{K}(S) \subseteq \bar{S}$: Let $f \in \mathcal{K}(S)$:

$$\begin{array}{ccc} x & \xrightarrow{s \in \bar{S}} & z \\ \downarrow f & \nearrow \sigma & \downarrow p \\ y & = & y \end{array} \rightsquigarrow \begin{array}{ccccc} x & = & x & = & x \\ \downarrow & & \downarrow s & & \downarrow p \\ y & \xrightarrow{\sigma} & z & \xrightarrow{id} & y \end{array}$$

Now, given

$$\Lambda = \{ \Lambda^n_j \hookrightarrow \Delta^n \}$$

We write

$$\mathcal{K}_R(\Lambda_i) = \{ \text{inner fibrations} \}$$

$$\mathcal{K}_L = \{ \text{left fibration} \}$$

$$\mathcal{K}_R = \{ \text{Right fibration} \}$$

Cor 1.3.16

Def: Anodyne $= \mathcal{K}_L \mathcal{K}_R(\Lambda_{i/L/\Gamma})$

Rem: f monomorphism $\leadsto f$ is anodyne iff it induces a weak eq. \circ realization.

Def: $\{ \text{Trivial fibration} \} = \mathcal{K}_R(\{ \partial \Delta^n \rightarrow \Delta^n \mid n \geq 0 \})$

$$\mathcal{K}_R(\partial \Delta) = \mathcal{K}_R(\bar{\partial} \Delta).$$

§2: Inner Fibrations

$$\text{So } \omega\text{-categories} = \mathcal{K}_R(\Lambda_i)$$

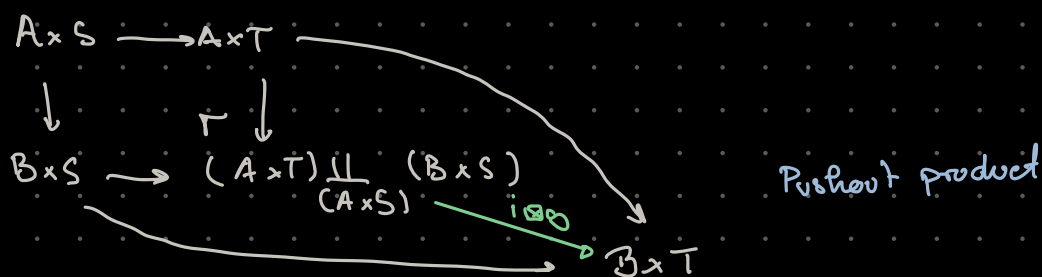
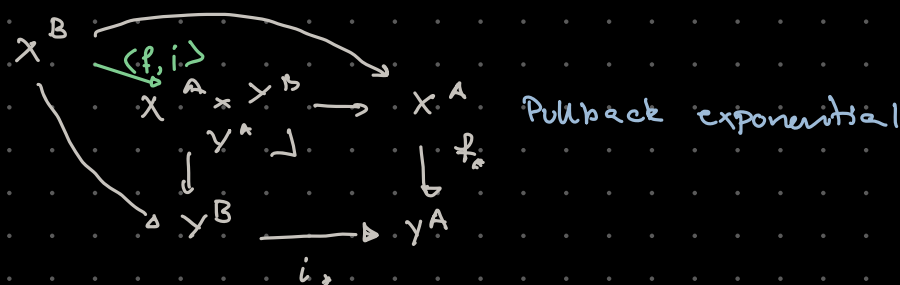
Recall: $C \in \text{Cat}_1$ is cartesian closed iff $b \in GC : C \rightarrow C$ admits a right adjoint, which we denote by $b \multimap b^a$

For sets ; $(\gamma^x)_n = \text{Hom}_{\text{sets}}(X \times \Delta^n, X)$

$$a \times c \rightarrow \rightarrow c \rightarrow^a$$

Consider:

$$A \hookrightarrow B, \quad \mathcal{B} \xrightarrow{\mathcal{B}} T, \quad X \xrightarrow{\mathcal{B}} Y$$



$$\sim \rightarrow i \boxtimes (-) \sim \langle -, i \rangle$$

$$\begin{array}{ccc} \text{Ar-Set} & \xrightleftharpoons[i \langle -, i \rangle]{i \boxtimes -} & \text{Ar-Set} \end{array}$$

Now, assume we are given the following diagram:

$$\begin{array}{ccc} S & \xrightarrow{\quad} & X^B \\ g \downarrow & \nearrow & \downarrow \\ T & \xrightarrow{\quad} & X^A \times_{Y^A} Y^B \end{array}$$

\cong

$$\begin{array}{ccc} S & \xrightarrow{\quad} & X^B \\ \downarrow & \nearrow & \downarrow \\ T & \xrightarrow{\quad} & X^A \times_{Y^A} Y^B \end{array}$$

these lifting problems are equivalent

$$\begin{array}{ccc} A \times S \rightarrow B \times S & & \\ \downarrow & \searrow & \downarrow \\ A \times T \rightarrow B \times T \rightarrow Y & & \end{array}$$

\cong

$$\begin{array}{ccc} (A \times T) \amalg (B \times S) & \xrightarrow{\quad} & X \\ (A \times S) \downarrow i \boxtimes g & \nearrow & \downarrow \\ B \times T & \xrightarrow{\quad} & Y \end{array}$$

Lem: 1.3.31 For i, g monomorphisms. Then $i \boxtimes g$ is respectively inner, L, R anodyne if either of i , or g resp., is.

Lem: $\Lambda = \{ \Lambda_j^n \hookrightarrow \Delta^n \mid n \geq 1, 0 < j < n \}$

$\Lambda = \text{anodyne maps} \supseteq \overline{\Lambda}$

$\Lambda^0 = \{ (K \hookrightarrow L) \boxtimes (\Lambda_j^2 \rightarrow \Delta^2) \mid K \hookrightarrow L \text{ mono} \}$

Prop: $X \xrightarrow{f} Y$ is an fib. of sSets, i monomorphisms, then $\langle f, i \rangle$ is resp. inner / L / R is f respectively is.

Moreover, if i is respectively inner / L / R anodyne, then $\langle f, i \rangle$ is a trivial fibration.

Pf:

$$\begin{array}{ccc} S \xrightarrow{\quad} X^B \\ g \downarrow \nearrow \exists? \downarrow \langle f, i \rangle & \Leftrightarrow & (A \times T) \amalg (B \times S) \xrightarrow{\quad} X \\ T \xrightarrow{\quad} X^A \times_{Y^A} Y^B & & \downarrow i \boxtimes g \nearrow \exists? \downarrow f \\ & & B \times T \xrightarrow{\quad} Y \end{array}$$

Ex: $X \xrightarrow{f} Y = \Delta^1, i \hookrightarrow B \rightsquigarrow$

$$\begin{array}{ccc} X^B & \xrightarrow{\quad} & X \\ \downarrow \langle f, i \rangle & \nearrow \text{inner} & \downarrow \\ \Delta^1 & \xrightarrow{\quad} & \Delta^1 \end{array}$$

Again ω -act

Def: Given $\mathcal{C}, \mathcal{D} \in \mathbf{Cat}_\omega$, set $\mathbf{Fun}(\mathcal{C}, \mathcal{D}) := \mathcal{D}^{\mathcal{C}} = \underline{\mathbf{Hom}}(\mathcal{C}, \mathcal{D}) \in \mathbf{Cat}_\omega$

Prop. $\mathcal{C} \in \mathbf{Set}$ is an ω -category iff $\mathcal{C}^{\Delta^2} \rightarrow \mathcal{C}^{\Delta^1}$ is a trivial fibration

Proof: Uses that

$$\overline{\Pi} = \{ (K \hookrightarrow L) \boxtimes (\Delta^2_1 \hookrightarrow \Delta^2) \}$$

Consider

$$\begin{array}{ccc} K & \xrightarrow{\quad} & \mathcal{C}^{\Delta^2} \\ \downarrow g & \nearrow & \downarrow \downarrow \\ L & \xrightarrow{\quad} & \mathcal{C}^{\Delta^2_1} \end{array} \qquad \begin{array}{ccc} K \times \Delta^2 \sqcup L \times \Delta^2 & \xrightarrow{\quad} & \mathcal{C} \\ \downarrow \downarrow & \nearrow & \downarrow \\ L \boxtimes \Delta^2_1 & \xrightarrow{\quad} & \Delta^0 \end{array}$$

Tommy Yang: Localization & coCartesian fibrations

26.10.2022 Lecture 4

Def: Let \mathcal{C} be an ∞ -category. A functor $f: \mathcal{C} \rightarrow \mathcal{C}[\mathcal{S}^{-1}]$ is a localization if

- 1) f takes \mathcal{S} to equivalences
- 2) $\text{Fun}(\mathcal{C}[\mathcal{S}^{-1}], \mathcal{D}) \simeq \text{Fun}^{\mathcal{S}}(\mathcal{C}, \mathcal{D})$ those functors which maps \mathcal{S} to equivalences.

Uniqueness almost follows by definition, so we'll only focus on existence.

Lem: $\Delta' \rightarrow \mathcal{J} = N(0 \rightrightarrows 1)$ is a localization

Pf: Note that elements in \mathcal{J}_k are of the form

$$j_0 \rightarrow \dots \rightarrow j_k \quad \text{w. } j_i = 0 \text{ or } 1$$

and non-degenerates of the form

$$0 \rightarrow 1 \rightarrow 0 \rightarrow 1 \rightarrow \dots$$

$$1 \rightarrow 0 \rightarrow 1 \rightarrow 0 \rightarrow \dots$$

$$\text{Fun}(\mathcal{J}, \mathcal{D}) \rightarrow \tilde{\text{Fun}}(\Delta', \mathcal{D}) \subset \text{Fun}(\Delta', \mathcal{D}).$$

Construct a filtration on \mathcal{J} : \mathcal{J}_k smallest sub-simplex that contains $\underbrace{0 \rightarrow 1 \rightarrow 0 \rightarrow 1 \rightarrow \dots}_{k+1}$. Consider the diagram

$$\Delta^k \rightarrow \mathcal{J}_{k-1}(\mathcal{J})$$

$$\downarrow \quad \quad \quad \downarrow$$

$$\Delta^k \rightarrow \mathcal{J}_k(\mathcal{J})$$

and

$$v_k: \Delta^k \rightarrow \mathcal{J}_k(\mathcal{J})$$

$$\Delta^k[n] \rightarrow \mathcal{J}_k(\mathcal{J})[n]$$

$$f: [n] \rightarrow [k] \mapsto f \pmod{2}$$

In particular: $f(j) = j+1 \mapsto \underbrace{1 \rightarrow 0 \rightarrow 1 \rightarrow \dots}_k$

$\mathcal{J}_{k-1}(\mathcal{J}) \rightarrow \mathcal{J}_k(\mathcal{J})$ is anodyne

$\leadsto \mathcal{F}_1(\mathcal{J}) = \Delta' \rightarrow \lim_n \mathcal{F}_n(\mathcal{J}) = \mathcal{J}$ is anodyne.

We wish to show that

$$\text{Fun}(\mathcal{J}, \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\Delta', \mathcal{D})$$

is an equivalence for all \mathcal{D} . We first note that this factors through

$$\text{Fun}(\mathcal{J}, \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\mathcal{F}_2(\mathcal{J}), \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\mathcal{F}_{2-1}(\mathcal{J}), \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\Delta', \mathcal{D})$$

Will show this is a trivial fibration

\leadsto show

$$\text{Fun}(\mathcal{J}, \mathcal{D}) \xrightarrow{\sim} \lim_n \text{Fun}^{\Delta'}(\mathcal{F}_n(\mathcal{J}), \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\Delta', \mathcal{D}) \xrightarrow{\sim} \text{Fun}^{\Delta'}(\mathcal{F}_1(\mathcal{J}), \mathcal{D})$$

is a trivial fibration

$$\begin{array}{ccccc} \partial\Delta^n & \rightarrow & \text{Fun}^{\Delta'}(\mathcal{F}_2(\mathcal{J}), \mathcal{D}) & \rightarrow & \text{Fun}^{\Delta'}(\Delta^{\mathbb{R}}, \mathcal{D}) \\ \downarrow & \nearrow & \downarrow & \nearrow & \downarrow \\ \Delta^0 & \xrightarrow{\sim} & \Delta^n & \rightarrow & \text{Fun}^{\Delta'}(\mathcal{F}_{2-1}(\mathcal{J}), \mathcal{D}) \rightarrow \text{Fun}^{\Delta'}(\Delta^{\mathbb{R}}, \mathcal{D}) \end{array}$$

wanna show \rightarrow lift exists, enough to show \rightarrow . By adj:

$$\begin{array}{ccccc} \Delta' & \rightarrow & \Delta^{\mathbb{R}} & \rightarrow & \text{Fun}(\Delta^n, \mathcal{D}) \xrightarrow{\text{ev}} \mathcal{D} \\ \downarrow & & \downarrow & & \downarrow \\ \Delta^{\mathbb{R}} & \rightarrow & \text{Fun}(\partial\Delta^n, \mathcal{D}) & \rightarrow & \mathcal{D} \end{array}$$

En equivalence

inner fibration

exists by Joyals

Lemma: $\forall \mathcal{C} \in \text{Cat}_{\infty} \exists$ localisation along \mathcal{C}_1 .

Pf:

$$\begin{array}{ccc} \coprod_{\text{ob}(\mathcal{C})} \Delta' & \rightarrow & \mathcal{C} \\ \downarrow & & \downarrow f \\ \coprod_{\text{ob}(\mathcal{C})} \mathcal{J} & \rightarrow & \mathcal{Y} \end{array}$$

Take an inner anodyne map $g: \mathcal{Y} \xrightarrow{\sim} \mathcal{X}$

$\Rightarrow g \circ f$ is anodyne.

Just some ∞ -cat.

Don't know a priori that this is an ∞ -category

Claim: \mathcal{X} is an ∞ -groupoid

Pf: Recall $h(-) \dashv N(-) \leadsto h(-)$ preserves pushouts, so applying it to the above diagram:

$$\begin{array}{ccc} \perp \{0 \rightarrow 1\} & \rightarrow & h\mathcal{C} \\ \downarrow & \lrcorner & \downarrow \\ \perp \{0 \rightrightarrows 1\} & \rightarrow & h\mathcal{Y} \end{array} \rightsquigarrow h\mathcal{Y} \text{ is a groupoid}$$

Claim: $\text{Fun}(X, \mathcal{D}) \rightarrow \text{Fun}^{\sim}(\mathcal{C}, \mathcal{D})$ is an equivalence.

Pf:

$$\begin{array}{ccccc} \text{Fun}(Y, \mathcal{D}) & \rightarrow & \text{Fun}^{\sim}(\mathcal{C}, \mathcal{D}) & \rightarrow & \text{Fun}(\mathcal{C}, \mathcal{D}) \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ \prod_{\mathcal{C}_1} \text{Fun}(\mathcal{I}, \mathcal{C}) & \rightarrow & \prod_{\mathcal{C}_1} \text{Fun}^{\sim}(\Delta', \mathcal{D}) & \rightarrow & \prod_{\mathcal{C}_1} \text{Fun}(\Delta', \mathcal{D}) \end{array}$$

\rightsquigarrow LHS is also a pullback

\rightsquigarrow Left lower horizontal map is a trivial fibration by above lemma.

\Rightarrow Same holds for upper horizontal map.

Hence we have showed that localization exists when considered along all morphisms, so now we wish to show existence w.r.t. any collection of morphisms.

Prop: For every $\mathcal{I} \subset \mathcal{C}_1$, \exists a localization along \mathcal{I} .

Pf: $\forall \mathcal{I} \subset \mathcal{C}_1 \exists$ a smallest sub ∞ -category $\mathcal{C}_{\mathcal{I}}$ containing \mathcal{I} :

$$\begin{array}{ccc} \mathcal{C}_{\mathcal{I}} & \rightarrow & \mathcal{C} \\ \downarrow & \lrcorner & \downarrow \\ \text{Localization of } \mathcal{C}_{\mathcal{I}} \text{ along all morphisms in } \mathcal{I} & \rightarrow & \mathcal{C} \end{array}$$

Again, we don't know if this is an ∞ -cat, so have to consider \mathcal{D} instead.

some ∞ -cat

Take an inner anodyne $\mathcal{W} \rightarrow \mathcal{D} \Rightarrow \mathcal{W} \rightarrow \mathcal{D}$ weak cat. eq. in Joyal's model

\rightsquigarrow

$$\begin{array}{ccccc} \text{Fun}(\mathcal{D}, \mathcal{E}) & \xrightarrow{g^*} & \text{Fun}(\mathcal{W}, \mathcal{E}) & \rightarrow & \text{Fun}^{\sim}(\mathcal{C}, \mathcal{E}) \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow \\ \text{Fun}(\mathcal{X}, \mathcal{E}) & \rightarrow & \text{Fun}^{\sim}(\mathcal{C}_{\mathcal{I}}, \mathcal{E}) & \rightarrow & \text{Fun}^{\sim}(\mathcal{C}, \mathcal{E}) \end{array}$$

Initial fibrations

$\rightsquigarrow g^*$ is an equivalence, thus $\text{Fun}(\mathcal{D}, \mathcal{E}) \simeq \text{Fun}^{\sim}(\mathcal{C}, \mathcal{E})$.

§ coCartesian fibrations

Motivation: For X a "nice" topological space, we have

$$\text{Cov}(X) \cong \text{Fun}(\Pi_{\leq 1}(X), \text{Set})$$

↖ covering spaces

Let $\mathcal{C} \in \text{Cat}_1$, $\mathbb{F}: \mathcal{C} \rightarrow \text{Cat}_1$, then we have the so called Grothendieck construction

$$\int_{\mathcal{C}} \mathbb{F} = \begin{cases} \text{obj: } (A, X), & A \in \mathcal{C}, X \in \mathbb{F}(A) \\ \text{morph: } (f, \varphi), & f: A \rightarrow B \in \text{Mor}(\mathcal{C}) \\ & \varphi: f_!(X) \rightarrow Y \in \text{Mor}(\mathbb{F}(B)) \\ & f_!: \mathbb{F}(A) \rightarrow \mathbb{F}(B) \text{ s.t. } \mathbb{F}(f) = f_! \end{cases}$$

Observation: $\int_{\mathcal{C}} \mathbb{F} \rightarrow \mathcal{C}$, $(A, X) \mapsto A$, can be characterized by functor that has enough coCartesian morphisms.

Def: Given $\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$, a map $\varphi: X \rightarrow Y$ is \mathcal{R} -coCartesian if there exists the following (unique) lift:

$$\left\{ \begin{array}{c} X \xrightarrow{\varphi} Z \\ \downarrow \varphi \\ \mathcal{D} \end{array} \right\} \xrightarrow{\quad} \mathcal{D} \xrightarrow{\quad} \mathcal{C}$$

$$\left\{ \begin{array}{c} \mathcal{R}(X) \xrightarrow{\mathcal{R}(\varphi)} \mathcal{R}(Z) \\ \mathcal{R}(\varphi) \downarrow \\ \mathcal{R}(X) \end{array} \right\} \xrightarrow{\quad} \mathcal{C}$$

(Note: The diagram shows a commutative square with a diagonal arrow from the top-left corner to the bottom-right corner, and a vertical arrow from the top-right corner to the bottom-right corner.)

$\mathcal{R}: \mathcal{D} \rightarrow \mathcal{C}$ is further called coCartesian if $\forall f: A \rightarrow B$ in \mathcal{C} and $X \in \mathcal{D}$ such that $\mathcal{R}(X) \rightarrow A$ there exists a \mathcal{R} -coCartesian morphism $\varphi: X \rightarrow Y$ s.t. $\mathcal{R}\varphi = f$.

Grothendieck correspondence: $\int_{\mathcal{C}}: \text{LFun}(\mathcal{C}, \text{Cat}) \xrightarrow{\cong} \text{coCart}(\mathcal{C})$

Time to generalise this to ∞ -categories: First idea is to take the nerve of the diagram used to define it in the 1-categorical case, which gives us

$$\begin{array}{ccc} \Lambda^2_0 & \rightarrow & N(\mathcal{D}) \\ \downarrow & & \downarrow \\ \Delta^2 & \rightarrow & N(\mathcal{C}) \end{array}$$

Def. Let $p: X \rightarrow \mathcal{S}$ be a functor of ssets. Then $f: \Delta \rightarrow X$ a morphism is said to be p -cartesian if $\forall n \geq 2$:

$$\begin{array}{ccccc} & & f & & \\ & \nearrow & & \searrow & \\ \Delta^{[0,1]} & \rightarrow & \Lambda^n_0 & \rightarrow & X \\ & \downarrow & \nearrow \text{Id} & \downarrow p & \\ & \Delta^n & \rightarrow & \mathcal{S} & \end{array}$$

We further say $p: X \rightarrow \mathcal{S}$ is coCartesian if p is an inner fibration, and there exists a lift

$$\begin{array}{ccc} \{0\} & \rightarrow & X \\ \downarrow \sigma & \nearrow & \downarrow p \\ \Delta' & \rightarrow & \mathcal{S} \end{array} \quad \text{s.t. } \sigma \text{ is } p\text{-cartesian in } X.$$

Ex. Right fibrations are Cartesian.

Next goal is to compare Cartesian fibrations with right fibrations.

Lem. A Cartesian fibration $p: \mathcal{E} \rightarrow \mathcal{C}$ is a right fibration \Leftrightarrow Every morphism in \mathcal{E} is p -Cartesian.

Prop. A Cartesian fibration $p: \mathcal{E} \rightarrow \mathcal{C}$ is a right fibration \Leftrightarrow for all $x \in \mathcal{C}$, the fibres \mathcal{E}_x are ∞ -groupoids.

Pf. \Rightarrow : \checkmark

\Leftarrow : We will show that every morphism in \mathcal{E} is p -Cartesian. Take any $f: \Delta' \rightarrow \mathcal{E}$, which takes z to y . We can then choose φ to be a p -Cartesian lift of $p(f)$ which has target y .

Consider the following diagram:

$$\begin{array}{ccccc} \Delta' & \xrightarrow{\varphi} & \Lambda^2_z & \rightarrow & \mathcal{E} \\ \downarrow & \nearrow \tau & \downarrow p & & \downarrow p \\ \Delta^2 & \rightarrow & \mathcal{C} & & \end{array} \quad \begin{array}{ccc} z & \xrightarrow{f} & y \\ z & \nearrow \varphi & \\ p(z) & \xrightarrow{p(f)} & p(y) \\ p(z) & \xrightarrow{\text{Id}} & p(z) \end{array}$$

The lift $\tau: \Delta^2 \rightarrow \mathcal{E}$ is given by

$$\begin{array}{ccc} z & \xrightarrow{f} & y \\ \psi \downarrow & \nearrow \varphi & \\ z & & \end{array}$$

∞ -groupoid

W. ψ a morphism in $\mathcal{E}p(z)$

$\Rightarrow \psi$ is an eq.

$\Rightarrow f$ is Cartesian.

□

Thomas Read: (Un)-Straightening

Lecture 5
02.11.22

$$\{\text{coCart fib over } \mathcal{C}\} \longleftrightarrow \{\text{functors } \mathcal{C} \rightarrow \text{Cat}_\infty\}$$

Intuitively, how does this work? Let $p: \mathcal{E} \rightarrow \mathcal{C}$ be a coCartesian fibration.

Given $x \in \mathcal{C} \leadsto \mathcal{E}_x \in \text{Cat}_\infty$. Given $f: x \rightarrow y$ we want $f_!: \mathcal{E}_x \rightarrow \mathcal{E}_y$.

Let $z \in \mathcal{E}_x$, take coCart lift of f

$$z \rightarrow f_! z$$

Given $\alpha: z \rightarrow z' \leadsto f_! \alpha$

$$\begin{array}{ccc} z & \xrightarrow{f_!} & f_! z \\ \alpha \downarrow & & \downarrow f_! \alpha \\ z' & \xrightarrow{f_!} & f_! z' \end{array}$$

because coCart
(This is unique up to const. choice)

In 1-Category land we would be done (the choices

would be unique), but in ∞ -category land this is more involved.

$$\begin{array}{ccc} z & \xrightarrow{f} & y \\ \parallel \alpha & & \parallel y \\ z & \xrightarrow{f} & y \end{array}$$

Def: KerSet , $f: K \rightarrow \mathcal{C}$, $p: \mathcal{E} \rightarrow \mathcal{C}$ coCart. Define $\text{Fun}_f^{\text{co}}(K, \mathcal{E})$ as the full subcategory spanned by those functors $K \rightarrow \mathcal{E}$ which lift f and sends all morphisms of K to coCartesian morphisms.

2-Simplices

Cor: Let $i: K \rightarrow L$ be left anodyne, $f: L \rightarrow \mathcal{C}$. Then

3.2.18

$$\text{Fun}_f^{\text{co}}(L, \mathcal{E}) \xrightarrow{i^*} \text{Fun}_{f \circ i}^{\text{co}}(K, \mathcal{E})$$

is a trivial fibration.

Recall that $\{0\} \hookrightarrow \Delta^n$ is left anodyne, so we get:

$$\text{Fun}_f^{\text{co}}(\Delta^1, \mathcal{E}) \xrightarrow{s} \text{Fun}_x^{\text{co}}(\{0\}, \mathcal{E}) \simeq \mathcal{E}_x$$

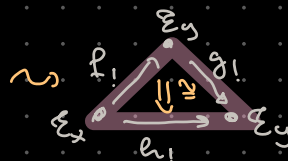
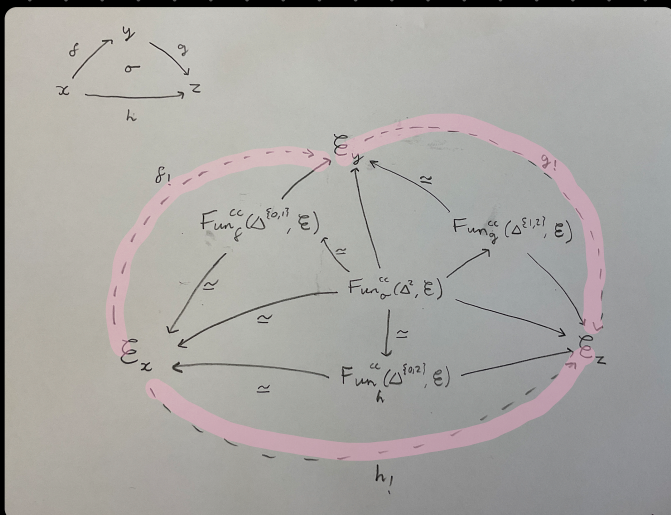
domain of f

is a trivial fibration, so it admits a section:

$$f_!: \mathcal{E}_x \xrightarrow{s} \text{Fun}_f^{\text{co}}(\Delta^1, \mathcal{E}) \rightarrow \mathcal{E}_y$$

An actual functor!

What does $\mathcal{C} \rightarrow \text{Cat}_\infty$ do on 2-simplices then?



This might get ugly, but defining functors between ∞ -categories are hard!

If A is a 1 category, then we have a functor $A \rightarrow \text{Cat}_{\infty}^1$ 1-cat of ∞ -cats

$$\leadsto N(A) \rightarrow N(\text{Cat}_{\infty}^1) \rightarrow N(\text{Cat}_{\infty}^1)[\text{Joy}^{-1}] \simeq \text{Cat}_{\infty}$$

Given $\omega \in A \leadsto N(A)[\omega^{-1}] \rightarrow \text{Cat}_{\infty}$

By choosing a clever choice of ω, A we can get $N(A)[\omega^{-1}] \simeq \mathcal{C}$.

Clever choice of A : $A = \Delta^{op}/\mathcal{C}$ category of simplices, defined by

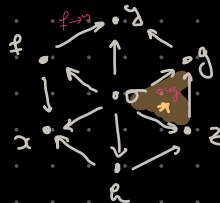
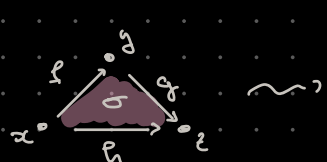
- Objects = $(n, \sigma: \Delta^n \rightarrow \mathcal{C})$

- Morphism $(n, \sigma) \rightarrow (m, \tau) = \alpha: [m] \rightarrow [n]$ s.t.

$$\begin{array}{ccc} \Delta^m & \xrightarrow{\alpha_*} & \Delta^n \\ \tau \searrow & \Downarrow \sigma & \swarrow \sigma \\ & \mathcal{C} & \end{array}$$

We see that $N(\Delta^{op}/\mathcal{C})$ is "barycentric subdivision".

$$\sigma: \Delta^2 \rightarrow \mathcal{C}$$



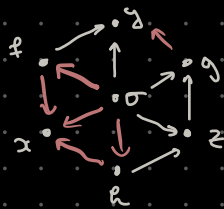
$$* = \sigma \rightarrow y \rightarrow z$$

Define $\mathcal{W}_{\mathcal{C}}$ to be the set of maps of the form

$$(n, \sigma) \rightarrow (m, \tau), \text{ with } \alpha: [m] \rightarrow [n],$$

such that $\alpha(0) = 0$.

So the red maps
are in $\mathcal{W}_\mathcal{C}$:



Want: $N(\Delta^{\text{op}}/\mathcal{C})[\mathcal{W}_\mathcal{C}] \rightarrow \mathcal{C}$ Joyal equivalence

Have canonical map

initial vertex map $\text{IV}: N(\Delta^{\text{op}}/\mathcal{C}) \rightarrow \mathcal{C}$

$$\begin{aligned} n\text{-simplex } [n_0] \xrightarrow{\alpha_0} [n_1] \rightarrow \dots \xrightarrow{\alpha_{n-1}} [n_n] &\mapsto \Delta^{\{0, \alpha_0(n_0), \alpha_1(\alpha_0(n_0)), \dots, \alpha_{n-1}(\alpha_0(n_0))\}} \\ \sigma: \Delta^{n_k} \rightarrow \mathcal{C} &\hookrightarrow \Delta^{n_k} \xrightarrow{\sigma} \mathcal{C} \end{aligned}$$

$w \in \mathcal{W}_\mathcal{C} \Rightarrow \text{IV}(w)$ degenerate, so IV induces $N(\Delta^{\text{op}}/\mathcal{C})[\mathcal{W}_\mathcal{C}] \rightarrow \mathcal{C}$.

First generalize to $X \in \text{Set}$, $W \subseteq X$

Define $L(X, W)$ as the pushout

$$\begin{array}{ccc} \coprod \Delta^i & \rightarrow & X \\ \downarrow & \searrow & \downarrow \\ \coprod \mathbb{J} & \rightarrow & L(X, W) \\ \downarrow \text{for } w_x & & \end{array}$$

from def of localization

\leadsto This satisfies that for $X = \mathcal{D}$ an ∞ -category:

$$L(\mathcal{D}, W) \simeq \mathcal{D}[W^{-1}] \quad \text{so it is indeed a generalization.}$$

So want to show that

$$L(N(\Delta^{\text{op}}/X), W_X) \rightarrow X$$

induced by IV

to be a weak categorical equivalence.

Fact: $X \mapsto L(N(\Delta^{\text{op}}/X), W_X)$ preserves monomorphisms and colimits.

\Rightarrow Suffices to consider the case $X = \Delta^n$:

$$\text{IV}: N(\Delta^{\text{op}}/\Delta^n) \rightarrow \Delta^n = N([n])$$

induced by

$$\widetilde{I\!V}: \Delta^p/\Delta^n \rightarrow [n]$$

$$(m, \sigma: \Delta^n \rightarrow \Delta^p) \mapsto \text{label of } \sigma(o)$$

Have

$$\begin{aligned} i: [n] &\longrightarrow \Delta^p \\ k &\longmapsto \Delta^{\{k, \dots, n\}} \hookrightarrow \Delta^p \end{aligned}$$

$$\widetilde{I\!V} \circ i \simeq \text{id}, \quad i \circ \widetilde{I\!V} \Rightarrow \text{id}$$

via morphisms in \mathcal{W}_{Δ^n}

So when taking the nerve we get a Joyal equivalence \simeq .

$\leadsto p: \mathcal{E} \rightarrow \mathcal{C}$ cartesian. Have

$$\Delta^p/\mathcal{C} \longrightarrow \text{Cat}_{\text{as}}$$

$$(n, \sigma: \Delta^n \rightarrow \mathcal{C}) \mapsto \text{Fun}_{\sigma}^{\text{cc}}(\Delta^n, \mathcal{E})$$

Given $(n, \sigma) \rightarrow (m, \tau)$ in $\mathcal{W}_{\mathcal{C}}$

$$\text{Fun}_{\sigma}^{\text{cc}}(\Delta^n, \mathcal{E}) \longrightarrow \text{Fun}_{\tau}^{\text{cc}}(\Delta^m, \mathcal{E})$$

$$\begin{array}{ccc} & \text{?} & \\ \sim \swarrow & & \searrow \sim \\ & \text{Fun}_{\text{col}}^{\text{cc}}(\Delta^o, \mathcal{E}) & \end{array}$$

By 2-out-of-3 this is again an eq.

So we get the straightening of $p: \mathcal{E} \rightarrow \mathcal{C}$.

$$\mathcal{C} \xrightarrow[\text{Joyal eq.}]{\text{Joyal}} N(\Delta^p/\mathcal{C})[\mathcal{W}_{\mathcal{C}}] \rightarrow \text{Cat}_{\text{as}}$$

Thm. Equivalence

$$\text{coCart}(\mathcal{C}) \simeq \text{Fun}(\mathcal{C}, \text{Cat}_{\text{as}})$$

$\xleftarrow{\text{Unstraightening}}$
 $\xrightarrow{\text{Straightening}}$

Special case:

- $\text{LFib}(\mathcal{C}) \simeq \text{Fun}(\mathcal{C}, \mathcal{S})$
- X a Kan complex: $\mathcal{S}/X \simeq \text{Fun}(X, \mathcal{S})$

A comment about unstraightening:

$$\text{id}: \text{Cat}_\infty \rightarrow \text{Cat}_\infty \quad \longleftrightarrow \quad p: (\text{Cat}_\infty)_{\times//} \rightarrow \text{Cat}_\infty \quad \text{universal coCart. fib.}$$

objects: pairs $(\mathcal{C}, x \in \mathcal{C}_0)$ ↗

mor $(\mathcal{C}, x) \rightarrow (\mathcal{D}, y)$ is $F: \mathcal{C} \rightarrow \mathcal{D}, \alpha: y \rightarrow Fx$

~,

$$\begin{array}{ccc} F^*(\text{Cat}_\infty)_{\times//} & \longrightarrow & \text{Cat}_{\times//} \\ F_p^* \downarrow & \lrcorner & \downarrow p \\ \mathcal{C} & \xrightarrow{F} & \text{Cat}_\infty \end{array}$$

Hannah MacDermott: Adjoint functor theorems

Talk 8 - 23.11.22

Cofinal functors

Def: $K, L, X \in \text{Set}$, $f: K \rightarrow L$, $p: X \rightarrow L$ inner fibration. We define the

$$\begin{array}{ccc} \text{Fun}_X(K, L) & \longrightarrow & \text{Fun}(K, X) \\ \downarrow & \lrcorner & \downarrow p_* \\ \Delta^0 & \xrightarrow{f} & \text{Fun}(K, L) \end{array}$$

$$\text{Fun}_L(L, X) = \text{Fun}_{\text{id}_L}(L, X)$$

$f: K \rightarrow L$ is cofinal if for all right fibrations $p: X \rightarrow L$, the map

$$\text{Fun}_L(L, X) \rightarrow \text{Fun}_f(K, X)$$

is a Joyal equivalence.

Thm: f is cofinal iff for any ∞ -category \mathcal{C} and $p: L \rightarrow \mathcal{C}$, the induced map

$$\mathcal{C}^{p/} \rightarrow \mathcal{C}^{p f/}$$

is a Joyal equivalence.

Cor: Joyal equivalences are cofinal.

Cor: p admits a colimit iff $p f$ does, and f preserves this colimit.

Thm: $f: \mathcal{C} \rightarrow \mathcal{D}$ map of simplicial sets, \mathcal{D} an ∞ -category. Then

$$f \text{ is cofinal} \Leftrightarrow \mathcal{C}_{d/} \text{ is weakly contractible for every object } d \text{ in } \mathcal{D}$$

$$|\mathcal{C}_{d/}| \simeq *$$

Prop: \mathcal{C}, \mathcal{D} ∞ -categories, $F: \mathcal{C} \rightarrow \mathcal{D}$. F admits a right adjoint iff \mathcal{C}/d admits a terminal object for each $d \in \mathcal{D}$.

Pf: Choose $d \rightsquigarrow (G_d, f: F G_d \rightarrow d)$ terminal object in \mathcal{C}/d . By 5.1.10, sufficient to prove that

$$\text{map}_{\mathcal{C}}(c, G.d) \rightarrow \text{map}_{\mathcal{D}}(F, FGd) \rightarrow \text{map}_{\mathcal{D}}(F, d)$$

is an equivalence. Consider the following diagram.

$$\begin{array}{ccccc} \text{map}_{\mathcal{C}/d}((c, \alpha), (Gd, \beta)) & \xrightarrow{\sim} & \text{map}_{\mathcal{D}/d}(c, \beta) & \xrightarrow{\sim} & \Delta^0 \\ \downarrow & & \downarrow & \textcircled{J} \text{ 3.3.19} & \downarrow \alpha \\ \text{map}_{\mathcal{C}}(c, G.d) & \xrightarrow{\sim} & \text{map}_{\mathcal{D}}(F, FGd) & \xrightarrow{\beta} & \text{map}_{\mathcal{D}}(F, d) \end{array}$$

\sim on components lying over the component of α

\leadsto True for all $\alpha: Fc \rightarrow d$.

Joyal equivalent to a small ω -cat. \square

Thm: 5.2.2. \mathcal{C} locally small cocomplete ω -category. If there exists $\mathcal{C}_0 \subseteq \mathcal{C}$ full essentially small

colimit dense, then

For any $c \in \mathcal{C}$, there is a

diagram s.t. $c = \text{colim}(K \rightarrow \mathcal{C}_0 \rightarrow \mathcal{C})$. $F: \mathcal{C} \rightarrow \mathcal{D}$ admits a right adjoint

\Leftrightarrow

F preserves colimits.

Def: Let S be a small set of objects in \mathcal{C} . Then S is weakly terminal if for every object $c \in \mathcal{C}$, there exists an object $s \in S$ s.t. $\text{map}_{\mathcal{C}}(c, s)$ is not empty.

An object t is said to be weakly terminal if the set $\{t\}$ is weakly terminal.

Lemma: 5.2.10. \mathcal{C} cocomplete, and there exists $\mathcal{C}_0 \subseteq \mathcal{C}$ essentially small

colimit dense full subcategory. Then \mathcal{C} has a weakly terminal object.

Pf: Consider $\mathcal{C}' \simeq \mathcal{C}$, $t = \text{colim}_{\mathcal{C}}(c' \rightarrow \mathcal{C}) = \text{colim}_{\mathcal{C}}(\mathcal{C}_0 \rightarrow \mathcal{C})$. Write c in \mathcal{C} as $\text{colim}_{\mathcal{C}}(K \rightarrow \mathcal{C}_0 \rightarrow \mathcal{C})$. Obtain a map $c \rightarrow t$.

Prop: 5.2.11. \mathcal{C} locally small cocomplete, $s \in \mathcal{C}$ weakly terminal and assume there exists a full subcategory $\mathcal{C}_0 \subseteq \mathcal{C}$ spanned by S . Then $\mathcal{C}_0 \rightarrow \mathcal{C}$ is cofinal

Pf: Sufficient to show that $(\mathcal{C}_0)_x$ is weakly contractible for all x in \mathcal{C}

\leadsto Sufficient to show that any functor

$$K \rightarrow (\mathcal{C}_0)_x$$

factors through $K * \Delta^0$ which is weakly cont.

\mathcal{C} cocomplete $\Rightarrow \mathcal{C}_x$ cocomplete 4.3.35

Take colimit cone

$$\begin{array}{ccc} K \rightarrow (\mathcal{C}_0)_{x_1} & \xrightarrow{\quad} & \mathcal{C}_{x_1} \\ \downarrow & \nearrow \mu & \\ K * \Delta^0 & & \end{array}$$

$\mu(\infty): x \rightarrow y$, pick $s \in S$ $y \rightarrow s$. Choosing a composite $x \rightarrow y \rightarrow s$ gives a 2-simplex $\Delta^2 \rightarrow \mathcal{C}$, adjoint to a map $\sigma: \Delta^1 \rightarrow \mathcal{C}^{x_1}$

$$\begin{array}{ccc} K * \Delta^0 \sqcup_{\Delta^0} \Delta^1 & \xrightarrow{\mu, \sigma} & \mathcal{C}_{x_1} \\ \downarrow & \nearrow \mu' & \\ K * \Delta^1 & & \end{array}$$

Lift exists because the vertical map is inner anodyne and \mathcal{C}_{x_1} is an ∞ -category. Restricting μ' to $K * \Delta^{1?}$ gives a map $\mu'': K * \Delta^1 \rightarrow \mathcal{C}_{x_1}$. $(\mathcal{C}_0)_{x_1}$ is full in \mathcal{C}_{x_1} , so sufficient to show that we can find $K * \Delta^1 \rightarrow \mathcal{C}_{x_1}$ which sends objects to $(\mathcal{C}_0)_{x_1}$.

$$\left. \begin{array}{l} \mu'' \text{ agrees w. } \mu \text{ on } K \\ \mu''(\infty) = \mu'(x_1) = \mu(x_1) = \sigma(x_1) = s \end{array} \right\} \text{Both in } \mathcal{C}_0$$

Cor: \mathcal{C} locally small cocomplete, $S \in \mathcal{C}$ weakly terminal. Then \mathcal{C} admits a terminal object.

Pf: Again take \mathcal{C}_0 to be the full subcategory spanned by S . \mathcal{C}_0 is small so $\mathcal{C}_0 \rightarrow \mathcal{C}$ has a colimit. $\mathcal{C}_0 \rightarrow \mathcal{C}$ is cofinal.

$$\leadsto \operatorname{colim}(\mathcal{C}_0 \rightarrow \mathcal{C} \xrightarrow{\text{id}} \mathcal{C}) = \operatorname{colim}(\mathcal{C} \xrightarrow{\text{id}} \mathcal{C})$$

Lemma 4.3.15

A colimit of the identity contains a terminal object.

Proof of adjoint functor theorem. Enough to show that every \mathcal{C}_{λ} has a terminal object

\leadsto Enough to show that $(\mathcal{C}_0)_{\lambda} \subseteq \mathcal{C}_{\lambda}$ essentially small, colimit dense full subcategory and \mathcal{C}_{λ} is locally small and cocomplete.

Need $\{\text{equivalence classes in } (\mathcal{C}_0)_{/d}\}$ to be small \sim but $\{\text{equivalence classes in } \mathcal{C}\}$

and $\text{map}(F_c, d)$ are all small.

Sufficient to prove that $(\mathcal{C}_0)_{/d} \in \mathcal{C}_{/d}$ is colimit-dense. Take $(c, \kappa: F_c \rightarrow d)$ in

$\mathcal{C}_{/d}$, and write

$$c = \text{colim}(K \rightarrow \mathcal{C}_0 \rightarrow \mathcal{C}).$$

Get a cone

$$K * \Delta^0 \rightarrow \mathcal{C}.$$

Want to lift this to $\mathcal{C}_{/d} \rightsquigarrow$ get a diagram

$$K \rightarrow \mathcal{C}_{/c} \xrightarrow{F} \mathcal{D}_{/F_c} \xrightarrow{\alpha} \mathcal{D}_{/d}$$

$$\begin{array}{ccc} K & \rightarrow & \mathcal{D}_{/d} \\ \downarrow & \nearrow & \uparrow \\ \mathcal{C}_0 & \xrightarrow{(\mathcal{C}_0)_{/d}} & \mathcal{C} \end{array}$$

$$\text{colim}(K \rightarrow \mathcal{C}_{0/d}) \text{ is } (c, \kappa) \quad \text{4.B.34}$$

MOTIVATION

There is useful analogies between homotopy theory and chain complexes with values in an abelian category (homological algebra)

- "homotopies" between chain maps
- "contractible" chain complexes

The analogue of the homotopy category of topological spaces is the 'derived category' of an abelian category, and analogue of stable homotopy theory is the homotopy category of spectra - Both which are triangulated categories, which gives a good setting for doing homological algebra. Working with triangulated categories can be very useful in practice but their theory can be very lacking, e.g:

- The category of functors between two triangulated categories doesn't inherit a triangulated structure
- Being triangulated is extra structure instead of a property
- Often can't use 'descent' ("gluing arguments")

Many of these problems can be traced back to the fact that we identify things without remembering why - they're identified. Using ∞ -categories lets us keep track of all this extra information.

We will therefore introduce the two ∞ -categorical analogues of these, see which properties these admit and use this as motivation for our definition.

Ex: \sum Let \mathcal{A} be a Grothendieck abelian category, i.e. locally presentable abelian category in which the small filtered colimit of a collection of SES's is again a SES.

Let $\text{Ch}(\mathcal{A}) = 1\text{-Cat. of chain complexes in } \mathcal{A}$

$$\text{Model structure on } \text{Ch}(\mathcal{A}) = \begin{cases} \text{cofibs} = \text{levelwise monomorphisms} \\ \text{weak eq.s} = \text{quasi-isomorphism} \end{cases}$$

We then define the (unbounded) derived ∞ -category of \mathcal{A} to be

$$\mathcal{D}(A) := \text{Nag}(\text{Ch}(A)^{\circ}),$$

the differential graded nerve of the bifibrant objects w.r.t. this model structure

(-) The localization of the category of chain complexes + the class of quasi-isomorphisms

Ex: First recall the ∞ -category of spaces $\text{N}^{\text{ch}}(\text{Kan}) = \mathcal{S}$. We wish to consider the "pointed" spaces:

$$\mathcal{S}_* \subseteq \text{Fun}(\Delta, \mathcal{S})$$

spanned by those functors which takes 0 to a "terminal" object of \mathcal{S} .

We define the ∞ -category of spectra as

$$\text{Sp} := \lim (\dots \xrightarrow{\Omega} \mathcal{S}_* \xrightarrow{\Omega} \mathcal{S}_* \xrightarrow{\Omega} \mathcal{S}_*)$$

↳ Ω is the loop functor which we will define later

↳ Think of it as a collection of spaces $\{X_n\}_{n \in \mathbb{Z}}$ + $X_n \overset{\text{empty}}{\cong} \Omega X_{n+1}$.

PROPERTIES THEY SHARE:

1) Pointed = Admits a zero-object

= An object which is both initial & terminal

- In Sp : spectrum which is a point at each level.
= Those whose homotopy groups all vanish
- In $\mathcal{D}(A)$: Acyclic complexes
= Those whose homology all vanish

Rem: Initial and terminal object is only unique up to contractible space of choices. We see this already at 1-categorical level: E.g. any one-point set is terminal in the category of sets, but $\{0\}$ and $\{1\}$ are two different (yet isomorphic) terminal objects.

Rem: In a pointed ∞ -category, we get a 'zero morphism' between any two objects by composing the initial/terminal map

$$x \rightarrow 0 \rightarrow y$$

2) Concept of (co)fibres or (co)kernels

- In $\mathcal{D}(\mathcal{A})$: Given a chain map $f: C \rightarrow D$ between chain complexes
 \leadsto Associated SES

$$0 \rightarrow C \rightarrow \text{cyl}(f) \rightarrow \text{cone}(f) \rightarrow 0$$

where $\text{cyl}(f)$ is a chain complex quasi-isomorphic to D .

\hookrightarrow We can often say interesting things about f by studying $\text{cone}(f)$, e.g. there is a LES relating $H_*(C)$, $H_*(D)$ and $H_*(\text{cone}(f))$, so $\text{cone}(f)$ is acyclic iff f is a quasi-isomorphism.

\hookrightarrow We can view $\text{cone}(f)$ as a 'cokernel' or 'cofibre' of f .

- In Sp : If we have a map $f: X \rightarrow Y$ between spectra, we get a fibre sequence

$$\text{fib}(f) \rightarrow X \xrightarrow{f} Y$$

which comes with an associated LES on homotopy groups. There is also a notion of cofibre sequences, and it turns out that fibre sequences are exactly the same as cofibre sequences, i.e.

$$\text{cofib}(\text{fib}(f) \rightarrow X) \simeq Y.$$

Let's make this notion more precise:

Def: A triangle in a pointed ∞ -category is a diagram $\Delta^{\times} \Delta \rightarrow \mathcal{C}$ of the form

$$\begin{array}{ccc} X & \rightarrow & Y \\ \downarrow & & \downarrow \\ 0 & \rightarrow & Z \end{array}$$

Rem: More explicitly: A triangle consists of

i) Two 1-simplices $X \xrightarrow{f} Y \xrightarrow{g} Z$

ii) A 2-simplex

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ & \searrow g & \downarrow h \\ & & Z \end{array}$$

witnessing a composite h of f and g .

iii) A 2-simplex

$$\begin{array}{ccc} X & & \\ \downarrow & \searrow h & \\ O & \rightarrow & Z \end{array}$$

witnessing a null-homotopy of h .

Rem: A $\Delta' \times \Delta' \rightarrow \mathcal{C}$ diagram $\begin{array}{ccc} X & \rightarrow & O \\ \downarrow & & \downarrow \\ Y & \rightarrow & Z \end{array}$ determines a ^{unique!} triangle by precomposing with

the map $\Delta' \times \Delta' \rightarrow \Delta' \times \Delta'$ that flips the two factors.

Def: Let \mathcal{C} be pointed and $p: \Delta' \times \Delta' \rightarrow \mathcal{C}$ the triangle

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow & & \downarrow g \\ O & \rightarrow & Z \end{array}$$

We say p is a fibre sequence if p is a pullback diagram (we say that p is a fibre of g). Dually, we say that p is a cofibre sequence if p is a pushout diagram (p is a cofibre of f).

Now, let's finally define stable:

Def: An ∞ -category \mathcal{C} is stable if

i) It is pointed

ii) Every morphism admits a fibre and a cofibre sequence

iii) A triangle is a fibre sequence iff it is a cofibre sequence

Rem: So being stable is a property, not extra structure.

↳ Think of given a set X , then it doesn't make sense to ask if it is a

group without also specifying $m: X \times X \rightarrow X$ group multiplication. On the other hand, it does make sense to ask if a group G is abelian.

↳ Why does this difference matter? Often easier to formulate cleaner abstract arguments when working w. properties.

Prop: [HA.1.4.2.27] Let \mathcal{C} be pointed. Then the following is equivalent,

- i) \mathcal{C} is stable
- ii) \mathcal{C} has finite colimits and the suspension functor $\Sigma: \mathcal{C} \rightarrow \mathcal{C}$, determined by the existence of the pushout square

$$\begin{array}{ccc} X & \rightarrow & 0 \\ \downarrow & \lrcorner & \downarrow \\ 0 & \rightarrow & \Sigma X, \end{array}$$

is an equivalence

- iii) \mathcal{C} has finite limits, and the loop functor $\Omega: \mathcal{C} \rightarrow \mathcal{C}$, determined by the pullback square

$$\begin{array}{ccc} \Omega X & \rightarrow & 0 \\ \downarrow & \lrcorner & \downarrow \\ 0 & \rightarrow & X, \end{array}$$

is an equivalence

Pf: We will show $i) \Leftrightarrow (ii)$, the case of (i) is dual. It is clear that $(i) \Rightarrow (ii)$, so assume \mathcal{C} admits finite colimits, and that Σ is an equivalence.

Claim 1: Every cofiber sequence is a fiber sequence

Pf: Let $p: \Delta' \times \Delta' \rightarrow \mathcal{C}$ denote the cofiber sequence determined by the square

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow & \lrcorner & \downarrow p \\ 0 & \rightarrow & C, \end{array} \quad (*)$$

and

$$p_0: \Delta' \amalg_{\Delta' \cap \Delta'} \Delta' \rightarrow \mathcal{C}$$

the restriction of p to the lower and right edges. To show the square is Cartesian, we need to show that the induced map

$$G/p \rightarrow G/p_0$$

is an equivalence (trivial Kan fibration). Sufficient to show that for each $X \in G$

$$G/p \times_G \{X\} \rightarrow G/p_0 \times_G \{X\}$$

is a weak equivalence of Kan complexes. We see that $G/p_0 \times_G \{X\}$ classifies diagrams of the form

$$\begin{array}{ccc} X & \rightarrow & B \rightarrow 0 \\ \downarrow & & \downarrow p \\ 0 & \rightarrow & C \rightarrow \Sigma A \end{array} \quad \begin{array}{l} \text{by taking the cofiber} \\ \text{Pushout} \end{array}$$

(Note that the cofiber of p is ΣA by $(*)$).

\leadsto The outer diagram

$$\begin{array}{ccccc} X & \rightarrow & B & \rightarrow & 0 \\ \downarrow & & \downarrow & \searrow & \downarrow \\ 0 & \rightarrow & C & \rightarrow & \Sigma A \end{array} \quad \begin{array}{l} \searrow \\ \Sigma X \end{array} \quad \Sigma X \rightarrow \Sigma A$$

$$\leadsto \quad \begin{array}{ccccc} \Sigma X & \searrow & & & \\ & \rightarrow & \Sigma A & \rightarrow & 0 \\ & \searrow & \downarrow p & & \downarrow \\ & & \Sigma B & \rightarrow & \Sigma C \end{array} \quad (\text{essentially uniquely extended!})$$

Let $q: \Delta' \times \Delta' \rightarrow G$ be the map corresponding to

$$\begin{array}{ccc} \Sigma A & \rightarrow & 0 \\ \downarrow p & & \downarrow \\ \Sigma B & \rightarrow & \Sigma C \end{array}$$

and $q_0 = q|_{\Delta' \times_{\Delta} \Delta'}$ be the restriction to the right and bottom edges. By the above we see that we get (a contractible choice of) maps of the form

$$G/p_0 \times_G \{X\} \rightarrow G/q \times_G \{\Sigma X\}$$

Hence we obtain a sequence of maps of the form

$$\mathcal{C}/p \times_{\mathcal{C}} \{x\} \rightarrow \mathcal{C}/p_0 \times_{\mathcal{C}} \{x\} \rightarrow \mathcal{C}/q \times_{\mathcal{C}} \{\Sigma x\} \rightarrow \mathcal{C}/q_0 \times_{\mathcal{C}} \{\Sigma x\}$$

These two compositions can be shown to agree with the maps induced by $\Sigma \rightsquigarrow$ weak equivalences!

2-out-of-6 \Rightarrow Each map is a weak equivalence!

Claim 2: Every map in \mathcal{C} has a fiber

Pf: Let $p: B \rightarrow C$ be a map and $A \in \mathcal{C}$ s.t. p admits a cofiber seq. of the form

$$\begin{array}{ccc} B & \rightarrow & 0 \\ p \downarrow & \lrcorner \downarrow & \\ C & \rightarrow & \Sigma A \end{array}$$

exists since ΣA is an eq.

From claim 1 we know this is also a fiber sequence:

$$\begin{array}{ccccc} A & \xrightarrow{\beta'} & B & \rightarrow & 0 \\ \downarrow \beta & & \downarrow p & \lrcorner \downarrow & \\ 0 & \rightarrow & C & \rightarrow & \Sigma A \end{array}$$

\leadsto Left hand square is again a pullback \checkmark

Claim 3: Every fiber sequence is a cofiber sequence

Pf: Consider \mathcal{C}^{op} . By claim II it has all cofibers. Since $\Sigma_{\mathcal{C}}$ is an equivalence, so is $\Sigma_{\mathcal{C}^{\text{op}}}$. Applying claim I to \mathcal{C}^{op} we get that all fiber sequences in \mathcal{C} is a cofiber sequence.

Cor: \mathcal{C} stable $\Rightarrow \mathcal{C}^{\text{op}}$ stable □

Cor: If \mathcal{C} is a stable, then Σ, Ω are inverse equivalences.

Pf: Let $X \in \mathcal{C}$, then we have a cofiber sequence

$$\begin{array}{ccc} X \rightarrow 0 & \text{Also fib seq} & X \hookrightarrow \Omega(\Sigma X) \rightarrow 0 \\ \downarrow \lrcorner \downarrow & \leadsto & \downarrow \lrcorner \downarrow \\ 0 \rightarrow \Sigma X & & 0 \longrightarrow \Sigma X \end{array}$$

□

Cor: $\Omega^{\infty} = \text{unique map} : S_p = \text{colim}(s_* \xrightarrow{\Omega} \dots \xrightarrow{\Omega} s_*) \rightarrow s_*$

$\Sigma^{\infty} : s_* \rightarrow S_p$

\leadsto Adjunction $S_p \overset{\Sigma^{\infty}}{\underset{\Omega^{\infty}}{\dashv}} s$

Def: A functor $f: \mathcal{C} \rightarrow \mathcal{D}$ between stable ∞ -categories are exact if it preserves the zero object (i.e. reduced) and (co)fiber sequences

Construction of mapping spectrum: We first note that

$$\text{Fun}^{\text{lex}}(\mathcal{C}, \mathcal{S}p) \xrightarrow[\Sigma^{\infty}_{\mathcal{C}(\cdot, -)}]{\text{preserves all limits}} \text{Fun}^{\text{lex}}(\mathcal{C}, \mathcal{S}_*)$$

Using that $\text{Map}_{\mathcal{C}}(X, -): \mathcal{C} \rightarrow \mathcal{S}$ preserves all limits, we get that it corresponds to

$$\text{map}_{\mathcal{C}}(X, -): \mathcal{C} \rightarrow \mathcal{S}p \quad \text{the space of } \text{map}_{\mathcal{C}}(X, Y) \text{ is}$$

$$\text{s.t.} \quad \text{Map}_{\mathcal{C}}(X, \Sigma^n Y)$$

$$\Omega^{\infty} \text{map}_{\mathcal{C}}(X, -) \simeq \text{Map}_{\mathcal{C}}(X, -).$$

All of this is natural in $X \in \mathcal{C}$, so can build the mapping spectrum

$$\text{map}_{\mathcal{C}}(-, -): \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{S}p.$$

Any stable ∞ -category is spectrally enriched.

TRIANGULATED STRUCTURE

As mentioned in the beginning, another approach to capturing the "stable" properties of the derived category and the category of spectra, is using triangulated categories.

As it turns out, the homotopy category of any stable ∞ -category can be equipped with a triangulated structure (actually 2).

First recall the following definitions:

Def: A 1-category is additive if

- 1) It is \mathcal{Ab} -enriched (Each hom-set carries the structure of an abelian group, and composition is bilinear)
- 2) Admits finite coproducts

Def: A triangulated category is an additive 1-category \mathcal{C} with an equivalence $\Sigma: \mathcal{C} \rightarrow \mathcal{C}$ together with a class of distinguished triangles, each which is a diagram of the form

$$X \rightarrow Y \rightarrow Z \rightarrow \Sigma X$$

satisfying the following:

TR1: i) $\forall X \in \mathcal{C}$, $X \xrightarrow{\text{id}} X \rightarrow 0 \rightarrow \Sigma X$ is a distinguished triangle

ii) $\forall u: X \rightarrow Y \exists Z \in \mathcal{C} + X \xrightarrow{u} Y \rightarrow Z \rightarrow \Sigma X$ distinguished

iii) Class of distinguished triangles are closed under iso.

TR2: $X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} \Sigma X$ distinguished iff

$Y \xrightarrow{g} Z \xrightarrow{h} \Sigma X \xrightarrow{-\Sigma f} \Sigma Y$ distinguished

TR3:

$$\begin{array}{ccccccc} X & \longrightarrow & Y & \longrightarrow & Z & \longrightarrow & \Sigma X \\ \downarrow u & & \downarrow & & \downarrow & & \downarrow \Sigma u \\ X' & \longrightarrow & Y' & \longrightarrow & Z' & \longrightarrow & \Sigma X' \end{array}$$

Distinguished

TR4: "Octahedral axiom"

The main point of writing out this definition is to appreciate how compact

The definition of stable ∞ -category is!

Thm: \mathcal{C} stable $\Rightarrow h\mathcal{C}$ admits a triangulated structure

We won't go into too many details, but we will describe why $h\mathcal{C}$ is additive and the distinguished triangles.

$h\mathcal{C}$ is additive: We define translation functor

$$[n]: \mathcal{C} \rightarrow \mathcal{C}$$

as n^{th} -fold suspension (for $n \geq 0$, $(-n)^{\text{th}}$ -fold for $n \leq 0$). Using that limits commutes w. $\text{Map}_{\mathcal{C}}(X, -)$ (and Ω is a limit), we get

$$\text{Map}_{\mathcal{C}}(X, Y) \simeq \text{Map}_{\mathcal{C}}(X, \Omega^n \Sigma^n Y) \simeq \Omega^n \text{Map}_{\mathcal{C}}(X, Y[n]) \text{ of ssets}$$

$\leadsto \text{Map}_{\mathcal{C}}(X, Y)$ is an ∞ -loop space, so $\pi_0 \text{Map}_{\mathcal{C}}(X, Y) \simeq \text{Hom}_{h\mathcal{C}}(X, Y)$ is an abelian group under loop sum

To see that $h\mathcal{C}$ admits finite coproducts, we see that \mathcal{C} admits such:

Note that

- 1) $X \simeq \text{cofib}(X[-1] \xrightarrow{u} 0)$
 - 2) $Y \simeq \text{cofib}(0 \xrightarrow{v} Y)$
 - 3) $u \perp v = (X[-1] \xrightarrow{0} Y)$ in $\text{Fun}(\Delta^1, \mathcal{C})$
 - 4) cofibers preserves colimits
- $\Rightarrow X \perp Y = \text{cofib}(X[-1] \xrightarrow{0} Y)$

Def: Let \mathcal{C} be a stable ∞ -category. We say that a diagram

$$X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} \Sigma X$$

in $h\mathcal{C}$ is a distinguished triangle, if there exists a diagram

$\Delta^1 \times \Delta^2 \rightarrow \mathcal{C}$ of the form

$$\begin{array}{ccccc} X & \xrightarrow{\tilde{f}} & Y & \rightarrow & 0 \\ \downarrow & \ulcorner & \downarrow \tilde{g} & \ulcorner & \downarrow \\ 0 & \rightarrow & Z & \xrightarrow{\tilde{h}} & \Sigma X \end{array}$$

s.t.

$$\begin{array}{ccc} \tilde{h} & \xrightarrow{\sim} & \omega \\ \downarrow \tilde{h} & \searrow \tilde{g} & \downarrow \tilde{u} \\ & \Sigma X & \end{array}$$

Remi: 1) Diagram 1 implies that $\omega \in \Sigma X$, which is the vertical map in the triangle

2) We made a choice of having triangles of the form $\Delta^1 \times \Delta^2 \rightarrow G$ rather than of the form $\Delta^2 \times \Delta^1 \rightarrow G$. Either choice works, but the two choices yields two different triangulated structures on $R\mathcal{G}$.

T-STRUCTURES

As mentioned in the beginning, a reason to consider stable ∞ -categories is that it is a form of analogue of chain complexes and gives us a framework to do (homological) algebra. An example of this is that given a filtered chain complex

$$\dots \subseteq F_{p-1} \subseteq F_p \subseteq \dots \subseteq C$$

we get a spectral sequence

$$E'_{p,q} = H_{p+q}(F_p/F_{p-1}) \Rightarrow H_{p+q}(C)$$

whose convergence is conditional on finiteness properties of the filtration. This spectral sequence arises from the SES of chain complexes

$$0 \rightarrow F_{p-1} \rightarrow F_p \rightarrow F_p/F_{p-1} \rightarrow 0.$$

This can be generalised to stable ∞ -categories by adding a so called "t-structure", which in particular associates an abelian category.

Notation: $\mathcal{C}[n] = \Sigma^n \mathcal{C}$, $\mathcal{C}[-n] = \Omega^n \mathcal{C}$

Def: Let \mathcal{C} be a stable ∞ -category. A t-structure on \mathcal{C} consists of a pair $(\mathcal{C}_{\geq 0}, \mathcal{C}_{\leq 0})$ of full subcategories of \mathcal{C} s.t.

- 1) $\mathcal{C}_{\geq 0}[1] \subseteq \mathcal{C}_{\geq 0}$ (closed under suspension) $(\mathcal{C}_{\geq 0} = \text{connective objects})$
 $\mathcal{C}_{\leq 0}[-1] \subseteq \mathcal{C}_{\leq 0}$ (closed under loop) $\mathcal{C}_{\leq 0} = \text{coconnective objects}$
- 2) $\forall X, Y \in \mathcal{C}: \text{Map}_{\mathcal{C}}(X, Y[1]) = 0$
- 3) $\forall X \in \mathcal{C} \exists$ fiber sequence

$$\begin{array}{ccc} \tau_{\geq 0} X & \rightarrow & X \\ \downarrow & & \downarrow \\ 0 & \rightarrow & \tau_{\leq -1} X \end{array} \quad \text{w.} \quad \begin{array}{l} \tau_{\geq 0} X \in \mathcal{C}_{\geq 0} \\ \tau_{\leq -1} X[1] \in \mathcal{C}_{\leq 0}. \end{array}$$

Write

$$\mathcal{C}_{\geq n} := \mathcal{C}_{\geq 0}[n], \quad \mathcal{C}_{\leq n} := \mathcal{C}_{\leq 0}[n]$$

Ex: ① $\mathcal{C} = \mathcal{C}h(R\text{-Mod}) \quad (D(R))$

$$\mathcal{C}_{\geq 0} = \{X \mid H_n X = 0, n < 0\}$$

$$\mathcal{C}_{\leq 0} = \{X \mid H_n X = 0, n > 0\}$$

② $\mathcal{C} = Sp$

$$\mathcal{C}_{\geq 0} = \{X \mid \pi_n X = 0, n < 0\}$$

$$\mathcal{C}_{\leq 0} = \{X \mid \pi_n X = 0, n > 0\}$$

Postnikov t-structure

Def: The inclusions of $\mathcal{C}_{\geq n}, \mathcal{C}_{\leq n} \hookrightarrow \mathcal{C}$ admits adjoints:

$$\begin{array}{ccccc} & \xrightarrow{\text{incl}} & & \xrightarrow{\tau_{\leq n}} & \\ \mathcal{C}_{\geq n} & \xleftarrow{\perp} & \mathcal{C} & \xleftarrow{\perp} & \mathcal{C}_{\leq n} \\ & \xleftarrow{\tau_{\geq n}} & & \xleftarrow{\text{incl}} & \end{array}$$

connective cover

\leadsto Fits into a commutative diagram (in sets)

$$\begin{array}{ccc} \mathcal{C}_{\geq n} & \xrightarrow[\text{incl}]{\tau_{\geq n}} & \mathcal{C} \\ \downarrow \tau_{\leq m} & \searrow & \downarrow \tau_{\leq m} \\ \mathcal{C}_{\geq n} \cap \mathcal{C}_{\leq m} & \xrightarrow[\tau_{\leq m}]{\text{incl}} & \mathcal{C}_{\leq m} \end{array}$$

$$\leadsto \partial: \tau_{\leq m} \circ \tau_{\geq n} \xrightarrow{\sim} \tau_{\geq n} \circ \tau_{\leq m}$$

of functors $\mathcal{C} \rightarrow \mathcal{C}_{\geq n} \cap \mathcal{C}_{\leq m}$.

Def: The heart of \mathcal{C} is $\mathcal{C}^\heartsuit := \mathcal{C}_{\geq 0} \cap \mathcal{C}_{\leq 0} \subseteq \mathcal{C}$

Rem: $\pi_n \text{Map}_{\mathcal{C}^\heartsuit}(X, Y) \cong \pi_n \text{Map}_{\mathcal{C}}(X, Y) \cong \pi_0 \text{Map}_{\mathcal{C}}(X, \Omega^n Y) = 0$ for $n > 0$
 since $\Omega^n Y \in \mathcal{C}_{\leq -n}$

$$\leadsto \mathcal{C}^\heartsuit \cong N(h(\mathcal{C}^\heartsuit)) \cong N(h(\mathcal{C}))^\heartsuit$$

Can be used to show it is abelian!

Ex: 1) $D(R)^\heartsuit = R\text{-Mod}$

2) $Sp^\heartsuit = Ah$

Def: $\pi_0 := \tau_{\leq 0} \circ \tau_{\geq 0} \cong \tau_{\geq 0} \circ \tau_{\leq 0}: \mathcal{C} \rightarrow \mathcal{C}^\heartsuit$

$$\pi_n: \mathcal{C} \xrightarrow{\Omega^n} \mathcal{C} \xrightarrow{\pi_0} \mathcal{C}^\heartsuit$$

Ex: 1) $\mathcal{C} = D(R) \Rightarrow \pi_n$ is the homology functor

2) $\mathcal{C} = Sp \Rightarrow \pi_n$ is the "normal" homotopy groups on spectra

Now, let's turn our focus to the spectral sequence promised as motivation for t -structures.

Def: A filtered object of $\mathcal{G} \in \text{Cat}_\infty$ is a functor $F: (\mathbb{Z}, \leq) \rightarrow \mathcal{G}$. We define the p^{th} graded piece of F as

$$gr_p(F) := \text{cofib}(F(p-1) \rightarrow F(p)).$$

From this we obtain an exact couple

$$\begin{array}{ccc} \pi_{p,q}(F(p)) & \xrightarrow{(1,-1)} & \pi_{p,q}(F(p)) \\ & \searrow (-1,0) & \downarrow (0,0) \\ & & \pi_{p,q}(gr_p(F)) \end{array}$$

Recall that π_i lands in the abelian 1-category \mathcal{G}^\heartsuit , where we can do all our normal homological algebra. So the exact couple gives us a spectral sequence

$$E'_{p,q} = \pi_{p,q}(gr_p(F))$$

in \mathcal{G}^\heartsuit with Serre differentials, i.e.

$$d^r: E'_{p,q} \rightarrow E'_{p-r, q+r-1}.$$

Under nice circumstances, this converges to $\pi_{p,q}(\text{colim } F)$ - e.g. if \mathcal{G} has sequential colimits (i.e. colimits for any diagram $(\mathbb{Z}_{\geq 0}, \leq) \rightarrow \mathcal{G}$), $\mathcal{G}_{\leq 0}$ is closed under sequential colimits and $F(p) \cong 0 \ \forall p < 0$.

Ex: $\mathcal{G} = \mathcal{D}(\mathbb{R}) \leadsto$ This SS corresponds to the classical spectral sequence on filtered colimits.

INTRODUCTION TO HIGHER ALGEBRA

§1: SYMMETRIC MONOIDAL \rightarrow OPERADS \rightarrow ∞ -LAND

The basic algebraic structure we want to generalise is

commutative monoid: Set M + multiplication $M \times M \rightarrow M$, unit $1 \in M$ s.t.

$$1x = x, xy = yx, x(yz) = (xy)z \quad \forall x, y, z \in M$$

\leadsto For categories this is a symmetric monoidal category: Category C + $1 \in C$ unit

When working with categories its unnatural to ask for $+ \otimes: C \times C \rightarrow C$

$X \otimes (Y \otimes Z) = (X \otimes Y) \otimes Z$, instead we want this structure to be given by extra data in form isomorphisms

$$\alpha_X: 1 \otimes X \simeq X$$

$$\beta_{X,Y}: X \otimes Y \simeq Y \otimes X$$

Just called monoidal if we do not assume this

$$\gamma_{X,Y,Z}: X \otimes (Y \otimes Z) \simeq (X \otimes Y) \otimes Z$$

+ coherence data!

It is clear that this approach can not be used to generalize to ∞ -categories. Lets spell out how to get an equivalent description of symmetric monoidal: The idea is that instead of giving the bifunctor $\otimes: C \times C \rightarrow C$, we instead for each n -tuple C_1, \dots, C_n and $d \in C$ we specify the collection of maps

$$C_1 \otimes \dots \otimes C_n \rightarrow d \quad + \text{ composition data (and coherence)}$$

This is done by colored operads (what I called multicategories in an earlier talk):

Def: A coloured operad \mathcal{O} consists of

- A set of objects $\text{ob } \mathcal{O}$ (sometimes called the colours)
- $\forall x_0, \dots, x_n, Y$ in \mathcal{O} : A set of multimorphisms $\text{Mul}_{\mathcal{O}}(x_0, \dots, x_n; Y)$
- Composition: Given

$$(x_1, \dots, x_{i_1}) \rightarrow y_1, \dots, (x_{i_{n-1}}, \dots, x_{i_n}) \rightarrow y_n$$

can compose this with $(y_1, \dots, y_n) \rightarrow z$ to obtain

$$(x_1, \dots, x_{i_n}) \rightarrow z$$

- Unit: $\text{id}_x: \{x\} \rightarrow x$

s.t. the composition law is unital and associative

Rem: Every colored operad \mathcal{O} has an underlying category by setting:

- objects = $\text{ob } \mathcal{O}$
- $\text{Hom}(X, Y) = \text{Mul}_{\mathcal{O}}(\{X\}, Y)$

~ So can view a coloured operad as a category + extra data in form of the collection of multimorphisms.

Rem: Given a (symmetric) monoidal category (C, \otimes) we can obtain a coloured operad C^\otimes w. underlying category C by assigning

$$\text{Mul}_C(X_0, \dots, X_n; Y) = \text{Hom}_C(X_0 \otimes \dots \otimes X_n, Y).$$

We can recover the symmetric monoidal structure of C^\otimes (up to canonical isomorphism) by Yoneda's Lemma. E.g. the tensor product $X \otimes Y$ is characterized by the fact that it corepresents the functor $Z \mapsto \text{Mul}_C(X, Y; Z)$.

~> Can consider symmetric monoidal categories as a special case of coloured operads.

To identify which assumptions it is on a coloured operad that makes it into a symmetric monoidal.

First recall:

Def: $\text{Fin}_* = \{\langle n \rangle\}$ pointed sets

$f: \langle n \rangle \rightarrow \langle m \rangle$ is inert if it takes some point to the basepoint and injective (isomorphic) on the rest.

$p: \langle n \rangle \rightarrow \langle 1 \rangle$ takes $j \mapsto 0$ if $j \neq i$, $i \mapsto i$, p is active if $p^{-1}(0) = \emptyset$

Construction: Let \mathcal{O} be a coloured operad ~> \mathcal{O}^\otimes category:

2.1.1.7.

- Objects = sequences of objects of \mathcal{O} X_0, \dots, X_n
- $\{X_i\}_{1 \leq i \leq m} \rightarrow \{Y_j\}_{1 \leq j \leq n}$ is given by a map $\alpha: \langle m \rangle \rightarrow \langle n \rangle$ in Fin_* together with a collection of multimorphisms

$$\{\varphi_j \in \text{Mul}_{\mathcal{O}}(\{X_i\}_{i \in \alpha^{-1}\{j\}}, Y_j)\}_{0 \leq j \leq n}$$

in \mathcal{O} .

- Composition of morphisms in \mathcal{O}^\otimes is determined by composition laws on Fin_* and on \mathcal{O} .

By construction, \mathcal{O}^\otimes comes equipped with a forgetful functor

$$\begin{aligned} \pi: \mathcal{O}^\otimes &\rightarrow \text{Fin}_* \\ (X_0, \dots, X_n) &\mapsto \langle n \rangle. \end{aligned}$$

Using π we can reconstruct the operad structure:

- Write $\mathcal{O}_{\langle n \rangle}^\otimes = \pi^{-1}\{\langle n \rangle\}$

- $p_i^1: \langle n \rangle \rightarrow \langle 1 \rangle \leadsto p_i^1: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}_{\langle 1 \rangle}^\bullet \cong \mathcal{O}$ - underlying category of \mathcal{O}
which induces equivalences: $\mathcal{O}_{\langle n \rangle}^\bullet \cong \mathcal{O}^{x_n}$.
 $(x_0, -, x_n) \mapsto \overline{x}$.

- $\text{Mul}_{\mathcal{O}}(x_0, -, x_n; Y) \Leftrightarrow \left\{ f: \overline{x} \rightarrow Y \text{ in } \mathcal{O}^\bullet \text{ s.t. } \pi(f): \langle n \rangle \rightarrow \langle 1 \rangle \text{ satisfies } \pi(f)^{-1}(0) = \{0\} \right\}$
 $\hookrightarrow \pi: \mathcal{O}^\bullet \rightarrow \text{Fin}_*$ determines $\text{Mul}_{\mathcal{O}}(x_0, -, x_n; Y)$ in the colored operad \mathcal{O}
- Can show that composition law for morphisms in \mathcal{O} can be recovered from the one in \mathcal{O}^\bullet .

\leadsto Can think of a coloured operad as an ordinary category \mathcal{O}^\bullet together with a forgetful functor $\pi: \mathcal{O}^\bullet \rightarrow \text{Fin}_*$ s.t. $\mathcal{O}_{\langle n \rangle}^\bullet \cong (\mathcal{O}_{\langle 1 \rangle}^\bullet)^{x_n}$.
It turns out that if π is further an "op-fibration" then this is exactly the symmetric monoidal categories.

So this is what we generalise to ∞ -land.

Def: An ∞ -operad is a functor of ∞ -categories $p: \mathcal{O}^\bullet \rightarrow N(\text{Fin}_*)$ s.t.

- 1) coCart. lift of every inert $f: \langle n \rangle \rightarrow \langle m \rangle$
 \hookrightarrow In particular it induces $f_!: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}_{\langle m \rangle}^\bullet$.
 - 2) For each $n \geq 0$, the functors $\{p_i^2: \mathcal{O}_{\langle n \rangle}^\bullet \rightarrow \mathcal{O}\}_{1 \leq i \leq n}$ determines an equivalence of categories $\mathcal{O}^{x_n} \cong \mathcal{O}_{\langle n \rangle}^\bullet$.
- + ...

So we in particular get that $(x_0, -, x_n) \in \mathcal{O}^{x_n}$ corresponds to an object of $\mathcal{O}_{\langle n \rangle}^\bullet$.

Def: A symmetric monoidal ∞ -category is an ∞ -operad $\mathcal{C}^\bullet \rightarrow \text{Fin}_*$ which is also coCartesian fibration.

Let's understand why this gives the desired structure:

Note that: $\langle n \rangle \rightarrow \langle m \rangle$ in $\text{Fin}_* \leadsto \mathcal{C}_{\langle n \rangle}^\bullet \rightarrow \mathcal{C}_{\langle m \rangle}^\bullet$

so the active morphisms $\langle 0 \rangle \rightarrow \langle 1 \rangle, \langle 2 \rangle \rightarrow \langle 1 \rangle$

\leadsto $\underbrace{\Delta^\bullet \rightarrow \mathcal{C}}_{\text{An object we denote by } 1}, \underbrace{\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}}_{\text{Our tensor}}$

The reason we went through all this work introducing ∞ -operads is because

they are exactly what we need to get nice algebraic structures, and symmetric monoidal categories is just one of many 'monoidal'-ish categories.

In general:

Def: We say that $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$ exhibits \mathcal{C} as \mathcal{O} -monoidal ω -category if

- 1) p is coCartesian
- 2) $\mathcal{O}^\otimes \xrightarrow{q} \text{Fin}_*$ is an ω -operad and $\mathcal{C}^\otimes \xrightarrow{p} \mathcal{O}^\otimes \xrightarrow{q} \text{Fin}_*$ exhibits \mathcal{C}^\otimes as an ω -operad.

Ex: Symmetric monoidal \leftrightarrow Fin_* -monoidal.

Rem: If \mathcal{C} is \mathcal{O} -monoidal by $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$, then any

$$X \in \mathcal{O}^\otimes_{\langle n \rangle} \leftrightarrow \{X_0, \dots, X_n\} \in \mathcal{O}$$

Given

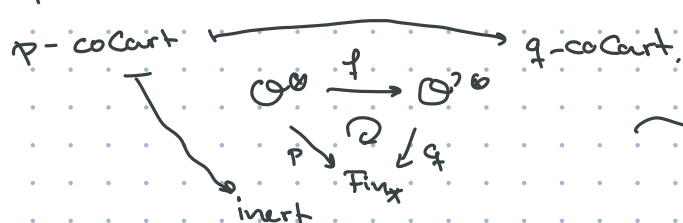
$$\{e \in \text{Mul}_{\mathcal{O}}(X_0, \dots, X_n; Y)\} \rightsquigarrow \prod_{0 \leq i \leq n} \mathcal{C}_{X_i} \xrightarrow{\mathcal{O}f} \mathcal{C}_Y$$

\uparrow
 fiber over x

It can be shown that one of the key elements is that p in particular is a morphism of ω -operads:

ALGEBRA OBJECTS

Def: A morphism of ω -operads :

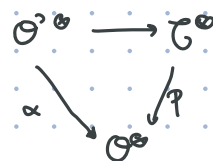


$$\text{Alg}_{\mathcal{O}}(\mathcal{O}') \subseteq \text{Fun}_{\text{Fin}_*}(\mathcal{O}^\otimes, \mathcal{O}'^\otimes)$$

i.e. morphism + fibration on underlinings

Def: Assume $p: \mathcal{C}^\otimes \rightarrow \mathcal{O}^\otimes$ be a fibration of ω -operads and given $\alpha: \mathcal{O}'^\otimes \rightarrow \mathcal{O}^\otimes$

$$\text{Alg}_{\mathcal{O}'/\mathcal{O}}(\mathcal{C}) \subseteq \text{Fun}_{\mathcal{O}^\otimes}(\mathcal{O}'^\otimes, \mathcal{C}^\otimes)$$



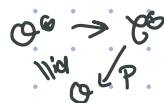
Equivalently:

$$\text{Alg}_{\mathcal{O}'/\mathcal{O}}(\mathcal{C}) = \text{fib. of } \text{Alg}_{\mathcal{O}}(\mathcal{C}) \xrightarrow{p} \text{Alg}_{\mathcal{O}}(\mathcal{O}) \text{ at } \alpha$$

$$\mathcal{O}'^\otimes \rightarrow \mathcal{C}^\otimes \mapsto \mathcal{O}'^\otimes \rightarrow \mathcal{C}^\otimes \xrightarrow{p} \mathcal{O}^\otimes$$

Case $\mathcal{O}^\otimes = \mathcal{O}'^\otimes, \alpha = \text{id}$:

$$\text{Alg}_{\mathcal{O}/\mathcal{O}}(\mathcal{C}) =: \text{Alg}_{/\mathcal{O}}(\mathcal{C})$$



Case $\Theta^{\otimes 0} = \Theta^{\otimes} = \text{Fin}_*$:

$$\text{CAlg}(\mathcal{C}) := \text{Alg}_{\Theta/\mathcal{C}}(\mathcal{C}) = \text{Alg}_{/\mathcal{C}}(\mathcal{C})$$

$$\begin{array}{ccc} \text{Fin}_* & \xrightarrow{\quad} & \mathcal{C}^{\otimes} \\ \parallel & \searrow & \\ \text{Fin}_* & & \end{array}$$

It's specific cases which gives us algebras!

Base case is associative algebra

Classically: An associative algebra object in a monoidal category \mathcal{C} is an object $A \in \mathcal{C}$ + unit map

$e: 1 \rightarrow A$ and multiplication $m: A \times A \rightarrow A$ s.t.

$$\begin{array}{ccc} 1 \otimes A & \xrightarrow{\text{ewid}} & A \otimes A \\ & \searrow & \swarrow m \\ & A & \end{array} \quad \begin{array}{ccc} A \otimes 1 & \xrightarrow{\quad} & A \otimes A \\ & \searrow & \swarrow m \\ & A & \end{array}$$

$$\begin{array}{ccc} A \otimes A \otimes A & \xrightarrow{\text{m} \circ \text{id}} & A \otimes A \\ \downarrow \text{id} \otimes m & & \downarrow m \\ A \otimes A & \xrightarrow{m} & A \end{array}$$

We will apply the formalism of ∞ -operads to introduce the notion of monoidal ∞ -category, and to each monoidal ∞ -category associate another ∞ -category $\text{Alg}(\mathcal{C})$ of associative algebra objects of \mathcal{C} .

Def: Colored operad **Assoc** (associative algebra)

- $\text{obAssoc} = \{a\}$
- $\text{Mul}_{\text{Assoc}}(\{a\}_{i \in I}, a) = \text{set of linear orderings on } I.$
any finite set.
- + composition...

\leadsto Obtain a category Assoc^{\otimes} by applying construction 2.1.1.7

Unwinding the def of Assoc^{\otimes} :

- $\text{obAssoc}^{\otimes} = \text{obFin}_*$
- $\langle m \rangle \rightarrow \langle n \rangle$ in $\text{Fin}_* \leadsto \langle m \rangle \rightarrow \langle n \rangle$ in Assoc^{\otimes} consists of a pair $(\kappa, \{i_1\}_{1 \leq i_1 \leq m})$ where $\kappa: \langle m \rangle \rightarrow \langle n \rangle$ is a map in Fin_* and \leq_{i_1} is a linear ordering on the inverse image $\kappa^{-1}\{i_1\} \subseteq \langle m \rangle$ for $1 \leq i_1 \leq n$

$\leadsto \text{Assoc}^{\otimes} = N(\text{Assoc}^{\otimes}) \in \text{Cat}_{\infty}$

Notation: We write $\text{Assoc} := \text{Assoc}^{\otimes} \times_{\text{Fin}_*} \{\langle 1 \rangle\}$

Note that as a simplicial set, Assoc is isomorphic to the 0-simplex Δ^0 ; But we use the notation Assoc to emphasize the role of the simplicial set as the underlying ∞ -category for the ∞ -operad Assoc^\otimes .

Def: A monoidal ∞ -category is a coCartesian fibration of ∞ -operads $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$.

Def: $\mathcal{C}^\otimes \in \text{Op}^\otimes$ equipped w. a fibration $q: \mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$. Then the ∞ -category of associative algebra objects of \mathcal{C} is $\text{Alg}(\mathcal{C}) := \text{Alg}_{/\text{Assoc}}(\mathcal{C})$ (∞ -operad sections of q)

$$\begin{array}{ccc} \text{Assoc}^\otimes & \xrightarrow{A} & \mathcal{C}^\otimes \\ \text{Id} \swarrow & & \searrow q \\ & \text{Assoc}^\otimes & \end{array}$$

Let's understand what this means: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a monoidal ∞ -category. Then

rem. 4.1.12-15 $\mathcal{C}_{\langle n \rangle}^\otimes \cong \mathcal{C}^\otimes \times_{\text{Assoc}^\otimes} \{\langle n \rangle\} \cong \mathcal{C}^n$, and for every linear ordering on $\{1, \dots, n\}$, the corresponding map $\langle n \rangle \rightarrow \langle 1 \rangle$ in Assoc^\otimes induces a functor $\mathcal{C}^n \rightarrow \mathcal{C}$. In particular

- $n=0 \leadsto$ unit object $1 \in \mathcal{C}$
- $n=2$, standard ordering on $\{1, 2\} \leadsto \otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$
- Evaluating on $a \in \text{Assoc}^\otimes$ determines a forgetful functor

$$\Theta: \text{Alg}(\mathcal{C}) \rightarrow \mathcal{C}$$

By abuse of notation we often identify $A \in \text{Alg}(\mathcal{C})$ with its image $\Theta(A)$ in \mathcal{C} . For each $n \geq 0$ a choice of ordering on $\{1, \dots, n\}$ determines an active morphism $\{a_i\}_{1 \leq i \leq n} \rightarrow a$ in Assoc^\otimes , which induces a morphism $\Theta(A)^{\otimes n} \rightarrow \Theta(A)$ in \mathcal{C} . In particular

- $n=2$ and standard ordering of $\{1, 2\}$: $m: \Theta(A) \otimes \Theta(A) \rightarrow \Theta(A)$
- \hookrightarrow Associative and unital up to homotopy

\leadsto In particular, it endows $\Theta(A)$ with the structure of an associative algebra object in $\text{h}\mathcal{C}$.

LEFT & RIGHT MODULES

Classically: \mathcal{C} monoidal category w. unit A , A associative algebra object of \mathcal{C} . A left A -module in \mathcal{C} is an object $M \in \mathcal{C}$ + action map $a: A \otimes M \rightarrow M$ s.t.

$$\begin{array}{ccc} A \otimes A \otimes M & \xrightarrow{m \circ id} & A \otimes M \\ \downarrow id \otimes a & & \downarrow a \\ A \otimes M & \xrightarrow{a} & M \end{array}$$

$$\begin{array}{ccc} 1 \otimes M & \xrightarrow{u \circ id} & A \otimes M \\ & \searrow & \swarrow a \\ & M & \end{array}$$

where $m = \text{multiplication}$
 $u = \text{Unit}$ } in A

\leadsto All left A -modules = $\text{LMod}_A(\mathcal{C})$

We wish to introduce a larger \mathcal{A} -operad $\text{LM}^{\mathcal{A}}$ which contains $\text{Assoc}^{\mathcal{A}}$. If $A \in \text{Alg}(\mathcal{C})$ we want a left A -module to be a map of \mathcal{A} -operads $M: \text{LM}^{\mathcal{A}} \rightarrow \mathcal{C}$ s.t. $M|_{\text{Assoc}^{\mathcal{A}}} = A$.

Def: Define a colored operad $\underline{\text{LM}}$ as:

- $\text{ob LM} = \{a, m\}$
- Let $\{X_i\}_{i \in I}$ be a finite collection of objects of LM . Then
 - $\hookrightarrow \text{Mul}_{\text{LM}}(\{X_i\}, a) = \begin{cases} \text{all linear orderings of } I \text{ if all } X_i = a \\ \text{Empty} & \text{otherwise} \end{cases}$
 - $\hookrightarrow \text{Mul}_{\text{LM}}(\{X_i\}, m) = \{ \text{All linear orderings } \{i_1, \dots, i_n\} \text{ of } I \text{ s.t. } X_{i_1} = m \text{ \& } X_{i_j} = a \text{ for } j > 1 \}$

Rem:
 4.2.1.2

$a \in \text{LM} \leadsto$ sub-colored operad of LM isomorphic to Assoc

a, a, \dots, a, m

Rem:

4.2.1.3

We first in this (1-categorical) case see how this can be used to understand modules.

Assume \mathcal{C} symmetric monoidal, $F: \text{LM} \rightarrow \mathcal{C}$ map of colored operads.

$\leadsto F|_{\text{Assoc}}: \text{Assoc} \rightarrow \mathcal{C} \leadsto$ Associative algebra object $F(a) = A \in \mathcal{C}$. Let $M = F(m) \in \mathcal{C}$.

Then the unique operation $\phi \in \text{Mul}_{\text{LM}}(\{a, m\}, m)$ determines $F(\phi): A \otimes M \rightarrow M$, which exhibits M as a left A -module.

\leadsto So $\begin{array}{l} \text{LM} \rightarrow \mathcal{C} \\ a \mapsto \text{associative algebra} \\ m \mapsto \text{left module over that algebra} \end{array}$

Notation: Apply construction 2.1.1.7 to LM to obtain the category $\text{LM}^{\mathcal{A}}$ from LM . Unwinding this construction we see that

$$1) \text{ ob LM}^{\mathcal{A}} = \{(\langle n \rangle, S) \mid S \subset \langle n \rangle^{\circ}\}$$

$$2) (\langle n \rangle, S) \rightarrow (\langle n' \rangle, S') = \alpha: \langle n \rangle \rightarrow \langle n' \rangle \text{ in } \text{Assoc}^{\mathcal{A}} \text{ s.t.}$$

$$i) S \cup \{*\} \xrightarrow{\alpha} S' \cup \{*\}$$

ii) $S' \in S' \Rightarrow \alpha^{-1}\{s'\}$ contains exactly one element of S , and that object is maximal w.r.t. the linear ordering of $\alpha^{-1}\{s'\}$.

Rem:

$$\begin{array}{ccc} LM & & LM^\otimes \\ a & \longleftarrow & (\langle 1 \rangle, \emptyset) \\ m & \longleftrightarrow & (\langle 1 \rangle, \langle 1 \rangle^\circ) \end{array}$$

Def: $LM^\otimes = N(LM^\otimes)$. This is an ∞ -operad via the forgetful map $LM^\otimes \rightarrow \text{Fin}_*$

Rem: The underlying ∞ -category LM of LM^\otimes is isomorphic to the discrete simplicial set $\Delta[1]$ w. two vertices, corresponding to $a, m \in LM$.

Rem: $\text{Assoc} \hookrightarrow LM \rightsquigarrow \text{Assoc}^\otimes \hookrightarrow LM^\otimes$, which is an isomorphism from Assoc^\otimes onto the full subcategory of LM^\otimes spanned by objects of the form $(\langle n \rangle, \emptyset)$.

Notation: Let $\mathcal{C}^\otimes \rightarrow LM^\otimes$ be a fibration of ∞ -operads. Write

$$\mathcal{C}_a^\otimes := \mathcal{C}^\otimes \times_{LM^\otimes} \text{Assoc}^\otimes, \quad \text{Underlying } \infty\text{-category } \mathcal{C}_a = \mathcal{C}^\otimes \times_{LM^\otimes} \{a\}$$

$$\mathcal{C}_m^\otimes = \mathcal{C}^\otimes \times_{LM^\otimes} \{m\}$$

Def: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a fibration of ∞ -operads, $q: \mathcal{C}^\otimes \rightarrow LM$ fibration of ∞ -operads s.t. $\mathcal{C}_a^\otimes \simeq \mathcal{C}^\otimes$. Write $\mathcal{C}_m^\otimes = M_0$ (Cat_∞) (Normally we say q exhibits M_0 as weakly enriched over \mathcal{C}^\otimes)

- $L\text{Mod}(M) := \text{Alg}_{/LM}(\mathcal{C})$ ∞ -category of left M -module objects of M_0 $\begin{array}{c} LM \rightarrow \mathcal{C} \\ \downarrow \\ LM \end{array}$
- Composition w. $\text{Assoc}^\otimes \hookrightarrow LM^\otimes$ determines a categorical fibration

$$L\text{Mod}(M_0) = \text{Alg}_{/LM}(\mathcal{C}) \xrightarrow{\circ \text{incl}} \text{Alg}_{/\text{Assoc}}(\mathcal{C}) = \text{Alg}(\mathcal{C})$$

$$\rightsquigarrow L\text{Mod}_A(M) = L\text{Mod}(M) \times_{\text{Alg}(\mathcal{C})} \{A\} \quad \infty\text{-category of left } A\text{-modules objects of } M_0.$$

\rightsquigarrow We think of $L\text{Mod}(\mathcal{C}_m)$ as given by pairs (A, M) where A is an associative algebra object of \mathcal{C}_a and M is a left A -module in \mathcal{C}_m .

Ex: $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ fibration of ∞ -operads, $\mathcal{C}^\otimes = \mathcal{C}^\otimes \times_{\text{Assoc}^\otimes} LM^\otimes$. Then " \mathcal{C}^\otimes exhibits \mathcal{C} as weakly enriched over \mathcal{C}^\otimes " \rightsquigarrow Can consider $L\text{Mod}(\mathcal{C}) = \text{Alg}_{LM/\text{Assoc}}(\mathcal{C})$

$$\begin{array}{ccc} LM^\otimes & \xrightarrow{\circ} & \mathcal{C}^\otimes \\ & \searrow & \downarrow \\ & & \text{Assoc}^\otimes \end{array}$$

Ex: Let $\mathcal{C}^\otimes \rightarrow \text{Assoc}^\otimes$ be a monoidal ∞ -category. Then

$$L\text{Mod}(\mathcal{C}) = \text{Alg}_{LM/\text{Assoc}}(\mathcal{C}) = LM^\otimes \xrightarrow{F} \mathcal{C}^\otimes$$

\downarrow
 Assoc^\otimes

$\rightsquigarrow F|_{\text{Assoc}} \in \text{Alg}(\mathcal{C})$, which we identify with its underlying object $F(a) = A \in \mathcal{C}$.

\rightsquigarrow Also have $F(m) = M \in \mathcal{C}$

The unique operation $q \in \text{Mod}_{LM}(\{a, m\}, m)$ determines a map

$$a: A \otimes M \rightarrow M \quad \text{in } \mathcal{C}$$

which is well-defined up to homotopy.

Since F is defined on all of \mathcal{LM}^\otimes , we get that this action map is compatible with the associative multiplication on A , up to coherent homotopy. In particular, if

$$\begin{aligned} m: A \otimes A &\rightarrow A && \text{multiplication on } A \\ u: 1 &\rightarrow A && \text{unit map} \end{aligned}$$

$$\begin{array}{ccc} A \otimes A \otimes M & \xrightarrow{m \otimes \text{id}} & A \otimes M \\ \downarrow \text{id} \otimes a & & \downarrow a \\ A \otimes M & \xrightarrow{a} & M \end{array} \quad \begin{array}{ccc} 1 \otimes M & \xrightarrow{u \otimes \text{id}} & A \otimes M \\ & \searrow & \swarrow a \\ & M & \end{array}$$

commutes up to homotopy. of ∞ -ops.

Rem: If $\mathcal{C}^\otimes \xrightarrow{q} \mathcal{LM}^\otimes$ is a cocartesian fibration, then the induced map

$\mathcal{C}_a^\otimes \rightarrow \text{Assoc}^\otimes$ is also a cocartesian fibration of ∞ -operads.

We further see that by straightening, q is classified by a map $\chi: \mathcal{LM}^\otimes \rightarrow \text{Cat}_\infty$, which is an " \mathcal{LM} -monoid object of Cat_∞ ". We have

$$\chi \in \text{Mon}_{\mathcal{LM}}(\text{Cat}_\infty) \simeq \text{Alg}_{\mathcal{LM}}(\text{Cat}_\infty) = \text{LMod}(\text{Cat}_\infty)$$

More informally: q can be thought of as giving an associative algebra \mathcal{C}_a in Cat_∞ together with a left module \mathcal{C}_m over \mathcal{C}_a . In particular q determines an action map

$$\otimes: \mathcal{C}_a \times \mathcal{C}_m \rightarrow \mathcal{C}_m$$

which is well-defined up to homotopy - and compatible with the monoidal structure on \mathcal{C}_a .